

Abstract:

A compact and efficient algorithm for the calculation of least squares splines is presented. Intended for use in interactive design of open and closed free-form curves, the algorithm performs parts of the computation in parallel to enable real-time curve-fitting. It employs the B-spline basis to form the normal equations and solves the normal equations by an envelope of $L D L^T$ factorization. A FORTRAN IV implementation is included.

A Real-Time Algorithm
for Least Squares Splines
and Its Application in
Computer-Aided Geometric Design

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1. Introduction

The creation and manipulation of free-form curves is one of the fundamental problems of computer aided geometric design. For objects such as ship hulls, automobile bodies, shoe lasts, or television tubes there are no canonical shapes: the form of the object is unrestricted. A design system must enable the designer to create, rapidly and accurately, the curves which describe these objects.

In many design systems, measurements from the designer's sketch specify the desired curve. From those measurements the design system should compute a compact and faithful approximation to that curve. Too often this approximation will include spurious oscillations not found in the data, yet may not reproduce cusps and inflection points that do exist in the data.

To overcome these difficulties, a combination of spline regression and B-spline Bezier curves may be employed [12]. The least squares spline provides a good initial approximation to the drawn curve and the Bezier polygon for this spline curve provides a convenient means for correcting defects in the approximation and for changing the curve. The theory and applications of Bezier curves are described elsewhere [8,10,12]; here the goal is an algorithm for the calculation of least squares splines in interactive

design systems.

In an interactive system the curve data are obtained from a graphics input device (typically a tablet) while the designer draws the curve. In such a system the least squares approximation to the curve should be displayed within seconds of the curve's completion, and if possible, that fit should be displayed while the curve is being drawn. In addition, the fitting program should be small enough for a local graphics processor.

This paper describes a least squares algorithm which meets these goals. Parts of the computation are overlapped to achieve fast response and to enable real-time curve-fitting. To minimize memory requirements only non-zero elements of the least squares equations are stored and the data points themselves are not stored at all. A well-documented FORTRAN IV implementation of the algorithm is included as Appendix I.

The remainder of the paper is divided into five sections. Sections 2 and 3 review, respectively, the properties of parametric spline curves and the development of the normal equations. In Section 4 the details of an efficient algorithm for the least squares spline calculation are described and the computational requirements for that algorithm are estimated. In Section 5, using the local convergence property of least squares splines, the algorithm

is re-structured as a pipeline, enabling real-time computation of the least squares spline and reducing storage requirements over the algorithm of Section 4. In the final section the algorithm is modified further for periodic curve fitting and the additional computational cost over the non-periodic curve fit is estimated.

2. Spline Curves

Polynomial splines are one of many possible generalizations of piecewise linear functions. Just as the piecewise linear spline is the function of least length connecting a set of points, so the cubic spline is the function of minimum strain energy connecting a set of points and having fixed slopes as its end points (Figure 2.1). From analogous minimum principles, higher order, odd degree polynomial splines and other, more general splines may be defined [17,18,20]. The cubic spline behaves much like the draftsman's spline -- a thin piece of bamboo held down by lead ducks -- for which it was named. Both cubic and draftsman's splines are widely used for approximating curves.

Alternatively, splines may be considered as piecewise polynomials -- a set of polynomials in the intervals of a partition joined so that derivatives match. For $k \geq 2$, the $(k-1)$ degree polynomial spline $S_k^\Delta(t)$ defined over the uniform knot set

$$(2.1) \quad \Delta: 0, h, \dots, m \cdot h, (m+1) \cdot h = a$$

is a polynomial of degree $(k-1)$ in each interval $(i \cdot h, (i+1) \cdot h)$ and has $(k-2)$ continuous derivatives over the complete interval $[0, a]$. This definition may be generalized further to allow non-uniform knots, but for efficiency and simplicity, the algorithms presented in this paper are

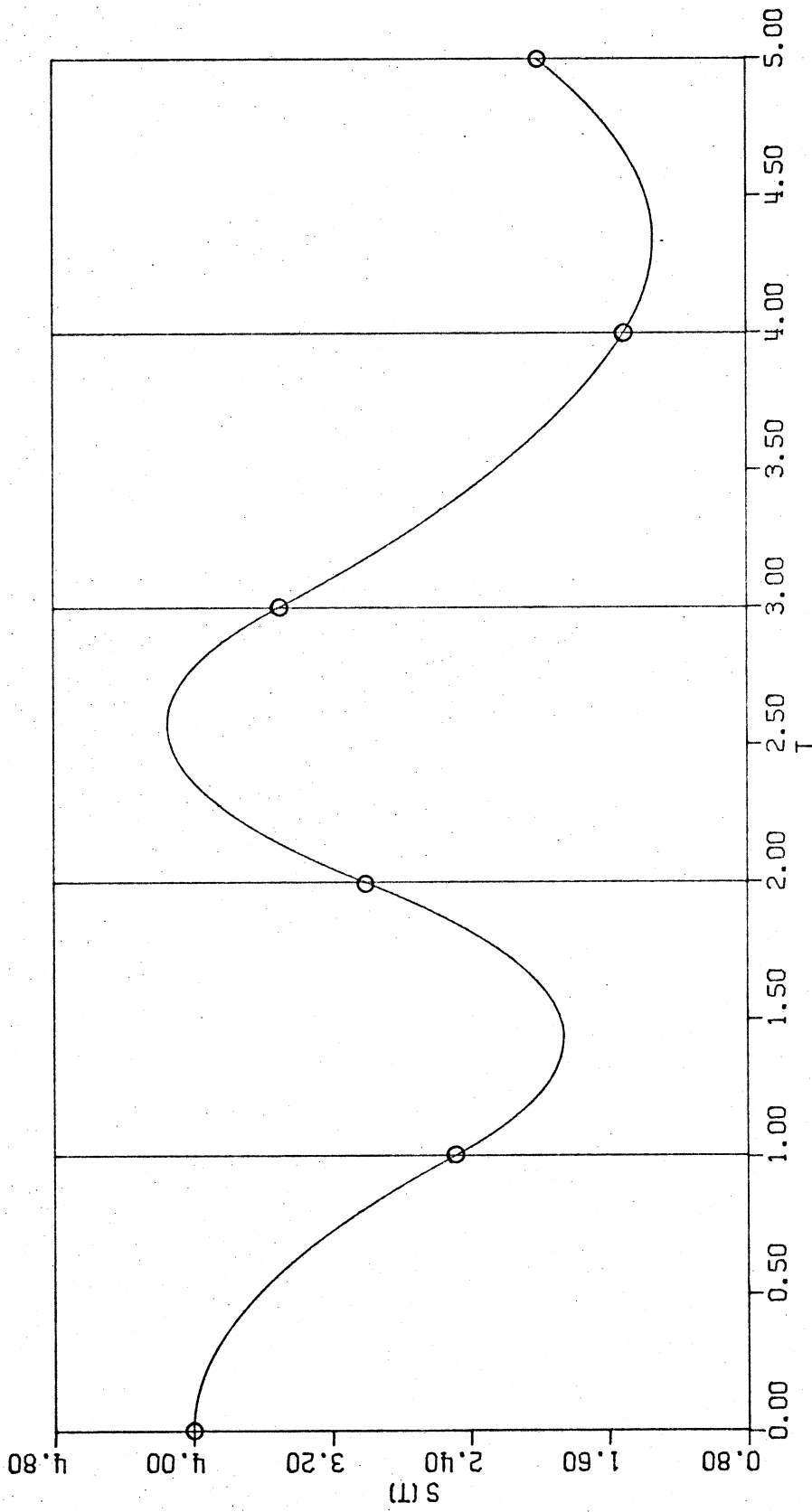


Figure 2.1: Cubic spline. The vertical lines indicate knot locations and the circles are the points which, with two end conditions, define the spline.

developed for uniform knots. The development for non-uniform knots is equally straightforward [6].

Since the spline is a set of simple polynomials, the spline may be evaluated from its polynomial coefficients

$$(2.2) \quad \{C_{i\ell}: 0 \leq i \leq m, 0 \leq \ell \leq k-1\},$$

where for t satisfying

$$(2.3) \quad i \cdot h \leq t < (i+1) \cdot h,$$

the value of the spline is given by

$$(2.4) \quad S(t) = \sum_{\ell=0}^{k-1} C_{i\ell} (t-ih)^\ell.$$

A third convenient characterization of a polynomial spline is in terms of the B-spline basis [1,2,3,4]. Any polynomial spline may be written as a linear combination

$$(2.5) \quad S(t) = \sum_{j=1}^{m+k} A_j N_{jk}(t)$$

of the B-spline basis functions N_{jk} (see Figure 2.2). The A_j are the $m+k$ B-spline basis coefficients.

The B-spline basis functions are themselves splines, and like any other spline, may be written as piecewise polynomials (see Appendix II). The B-splines are non-negative,

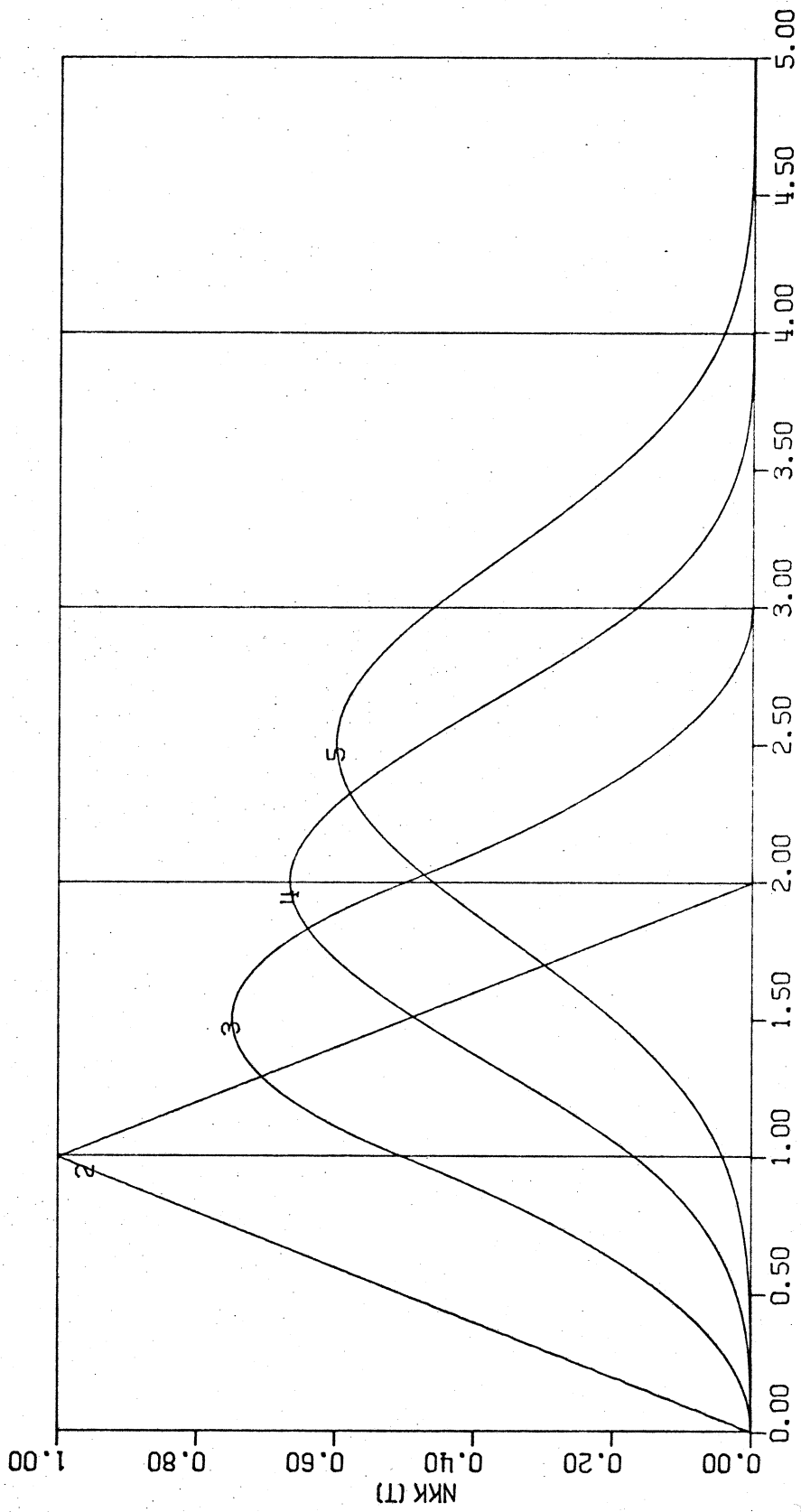


Figure 2.2: B-splines for $k = 2, 3, 4, 4, 5$. The vertical lines indicate knot locations.

$$(2.6) \quad N_{jk}(t) \geq 0 \text{ for } 0 \leq t \leq a,$$

normalized,

$$(2.7) \quad 1 = \sum_{j=1}^{m+k} N_{jk}(t) \text{ for } 0 \leq t \leq a,$$

and have local support,

for $0 \leq t \leq a$ and $1 \leq j \leq m+k$

$$(2.8) \quad N_{jk}(t) \neq 0 \text{ if and only if } (j-k) \cdot h < t < j \cdot h.$$

Since the B-spline basis is local, for

$$(2.9) \quad i \cdot h \leq t < (i+1) \cdot h$$

the sum (2.5) reduces to

$$(2.10) \quad S(t) = \sum_{j=i+1}^{i+k} A_j N_{jk}(t).$$

As an example of the three spline definitions, consider the following three different characterizations of a simple piecewise linear spline (Figure 2.1).

1) The function of least arc length connecting the points $(0,2)$, $(1,5)$, and $(2,1)$.

2) The piecewise polynomial

$$(2.11) \quad S(t) = \begin{cases} 2 + t \cdot 3 & \text{for } 0 \leq t < 1 \\ 5 + (1-t) \cdot 4 & \text{for } 1 \leq t \leq 2 \end{cases}.$$

3) The linear combination of B-spline basis functions

$$(2.12) S(t) = 2 \cdot N_{12}(t) + 5 \cdot N_{22}(t) + 1 \cdot N_{32}(t)$$

where

$$(2.13) N_{12}(t) = \begin{cases} 1-t & \text{for } 0 \leq t \leq 1 \\ 0 & \text{otherwise} \end{cases}$$

$$(2.14) N_{22}(t) = \begin{cases} t & \text{for } 0 \leq t < 1 \\ 2-t & \text{for } 1 \leq t < 2 \\ 0 & \text{otherwise} \end{cases}$$

$$(2.15) N_{32}(t) = \begin{cases} t-1 & \text{for } 1 \leq t \leq 2 \\ 0 & \text{otherwise} \end{cases}.$$

So far the discussion has been restricted to spline functions of one variable. "Spline curves" in two or more dimensions may be represented as two or more spline functions of a parameter t . In two dimensions, a parametric spline curve is given by

$$(2.16) \underline{S}(t) = (S_x(t), S_y(t)) \quad \text{for } 0 \leq t \leq a,$$

where $S_x(t)$ and $S_y(t)$ are one-dimensional splines.

For t satisfying

$$(2.17) i \cdot h \leq t < (i+1) \cdot h$$

the parametric spline may be written as a polynomial with vector coefficients

$$(2.18) \underline{S}(t) = \sum_{\ell=0}^{k-1} \underline{C}_{i\ell} (t - i \cdot h)^\ell$$

or as a linear combination of B-spline basis functions

$$(2.19) \quad \tilde{S}(t) = \sum_{j=i+1}^{i+k} \tilde{A}_j N_{jk}(t)$$

where the \tilde{A}_i are the vector B-spline basis coefficients.

3. Least Squares

Given measurements from the designer's drawn curve, the proposed design system must produce a spline curve which is in some sense "close" to the original. Furthermore, the resulting spline curve should have relatively few parameters so that it may be stored compactly and so that the Bezier polygon for the curve will have a small number of vertices [15].

One means of computing such a curve is least squares. Unlike an interpolate, the least squares spline need not pass through every data point; therefore it may have far fewer parameters than there are data. It will tend to smooth measurement noise, to flatten spurious "wiggles," and to be insensitive to gaps in the data [11]. In addition, as the following sections show, the least squares spline may be computed quite efficiently.

In one dimension the computation of the least squares spline is particularly simple. Given a set of weighted data

$$(3.1) \quad X = \{(t_q, y_q, w_q) : 1 \leq q \leq M\},$$

with abscissa t_q , ordinate y_q , and positive weight w_q , the least squares spline $S_k^\Delta(t)$ is that spline closest to the data in that it minimizes the weighted least squares distance

$$(3.2) \quad d(S, X) = \sum_{q=1}^M w_q \cdot (S(t_q) - y_q)^2$$

over all splines of order k and knot set Δ . For simplicity in the following development, the weights w_q are taken to be unity.

The traditional approach to this approximation problem is to form and solve the normal equations. The spline is written as a linear combination of some set of basis functions (here chosen as the B-splines)

$$(3.3) \quad S(t) = \sum_{j=1}^{m+k} A_j N_{jk}(t),$$

so that the distance functional $d(S, X)$ may be written as a quadratic function of the basis coefficients \underline{A}

$$(3.4) \quad F(\underline{A}) = d(S, X) = \underline{A}^T \cdot \underline{G} \cdot \underline{A} - 2 \cdot \underline{A}^T \cdot \underline{B} + C,$$

where the discrete Gram matrix G is given by

$$(3.5) \quad G = [g_{ij}] = \left[\sum_{q=1}^M N_{ik}(t_q) \cdot N_{jk}(t_q) \right],$$

the vector \underline{B} is given by

$$(3.6) \quad \underline{B} = [b_i] = \left[\sum_{q=1}^M y_q \cdot N_{ik}(t_q) \right],$$

and the constant C is given by

$$(3.7) \quad C = \sum_{q=1}^M y_q^2.$$

The minimization of F with respect to \underline{A} is a simple multivariate calculus problem. The unique minimum of F is attained at \underline{A}^* if the gradient

$$(3.8) \quad \frac{\partial F}{\partial \underline{A}} = 2 \cdot (G \cdot \underline{A} - \underline{B})$$

vanishes and the Hessian matrix

$$(3.9) \quad \left[\frac{\partial^2 F}{\partial A_i \partial A_j} \right] = 2 \cdot G$$

is positive definite. Therefore the matrix form of the normal equations

$$(3.10) \quad G \cdot \underline{A}^* = \underline{B}$$

must be satisfied at the minimum \underline{A}^* .

For typical graphics tablet data, and for any other data numerous and reasonably distributed with respect to the knots, the Gram matrix G will be positive definite and a unique minimum will exist [2,11]. Furthermore, under appropriate conditions, the matrix G will be well conditioned and the numerical stability of the method is assured [1,5, 11,16].

For least squares curves in two or more dimensions, a

different distance functional is minimized. This functional is constructed so that an n-dimensional problem reduces to n one-dimensional problems. The following development in two dimensions generalizes easily to curves in higher dimensions.

Given a set of data

$$(3.11) Y = \{(x_q, y_q) : 1 \leq q \leq M\},$$

the least squares spline

$$(3.12) S_k^\Delta(t) = \sum_{j=1}^{m+k} (A_j^x, A_j^y) \cdot N_{jk}(t)$$

is that spline which minimizes the distance functional

$$(3.13) d_2(S, Y) = \sum_{q=1}^M [(S_x(t_q) - x_q)^2 + (S_y(t_q) - y_q)^2]$$

over all parametric splines of given order k and knots Δ .

The parameterization of the data t_q is chosen as a piecewise linear approximation to the arc length

$$(3.14) \begin{aligned} t_1 &= 0 \\ t_q &= t_{q-1} + [(x_q - x_{q-1})^2 + (y_q - y_{q-1})^2]^{1/2} \\ &\text{for } 2 \leq q \leq M. \end{aligned}$$

As before, the distance functional $d_2(S, Y)$ may be written as a

quadratic function of \tilde{A}^x and \tilde{A}^y

$$(3.15) \quad F(\tilde{A}^x, \tilde{A}^y) = \begin{aligned} & (\tilde{A}^x)^T \cdot G \cdot \tilde{A}^x - 2 \cdot (\tilde{A}^x)^T \cdot \tilde{B}^x + C^x + \\ & (\tilde{A}^y)^T \cdot G \cdot \tilde{A}^y - 2 \cdot (\tilde{A}^y)^T \cdot \tilde{B}^y + C^y, \end{aligned}$$

where the discrete Gram matrix G is given by

$$(3.16) \quad G = [g_{ij}] = \left[\sum_{q=1}^M N_{ik}(t_q) \cdot N_{jk}(t_q) \right],$$

the vectors \tilde{B}^x and \tilde{B}^y are given by

$$(3.17) \quad \tilde{B}^x = [b_i^x] = \left[\sum_{q=1}^M x_q \cdot N_{ik}(t_q) \right]$$

$$(3.18) \quad \tilde{B}^y = [b_i^y] = \left[\sum_{q=1}^M y_q \cdot N_{ik}(t_q) \right],$$

and the constants C^x and C^y are given by

$$(3.19) \quad C^x = \sum_{q=1}^M x_q^2$$

$$(3.20) \quad C^y = \sum_{q=1}^M y_q^2.$$

At the least squares solution (A_x^*, A_y^*) the gradients

$$(3.21) \quad \frac{\partial F}{\partial \tilde{A}^x} = 2 \cdot (G \cdot \tilde{A}^x - \tilde{B}^x)$$

$$(3.22) \quad \frac{\partial F}{\partial \tilde{A}^y} = 2 \cdot (G \cdot \tilde{A}^y - \tilde{B}^y)$$

must vanish so that the normal equations

$$(3.23) \quad G \cdot \tilde{A}^X = \tilde{B}^X$$

$$(3.24) \quad G \cdot \tilde{A}^Y = \tilde{B}^Y$$

must be satisfied.

The two equations (3.23, 3.24) are identical to the one-dimensional normal equation (3.10) and may be solved separately by the techniques used for the one-dimensional problem. Note that the same matrix G appears in both equations.

4. Computational Considerations

The spline curve-fitting method of Section 3 may be implemented quite efficiently. In this section a detailed implementation of the calculation is presented and the work required is estimated. These work estimates will be used in Section 5 to determine the relative time required for various parts of the calculation.

The first problem is the evaluation of the spline itself. A polynomial spline is most efficiently evaluated from its piecewise polynomial representation (2.2-4). Given a spline of order k , with knots

$$(4.1) \quad \Delta: 0, h, \dots, m \cdot h, (m+1) \cdot h = a$$

and piecewise polynomial coefficients

$$(4.2) \quad \{C_{i\ell}: 0 \leq i \leq m, 0 \leq \ell \leq k-1\},$$

the spline can be computed by Horner's rule

$$(4.3) \quad \begin{aligned} S_0 &= C_{i,k-1} \\ S_\ell &= (t - i \cdot h) \cdot S_{\ell-1} + C_{i,k-1-\ell} \quad \text{for } 1 \leq \ell \leq k-1 \\ S(t) &= S_{k-1} \end{aligned}$$

for

$$(4.4) \quad i \cdot h \leq t < (i+1) \cdot h.$$

For the cubic spline ($k = 4$) with $\delta = t - i \cdot h$, the preceding is simply

$$(4.5) \quad S_4(t) = C_{i0} + \delta \cdot (C_{i1} + \delta \cdot (C_{i2} + \delta \cdot C_{i3})),$$

requiring three multiplications. In general, the evaluation of a spline of order k requires $(k-1)$ multiplications.

Using this method to evaluate the basis functions, the normal equations are computed. The Gram matrix G of the normal equations is symmetric by definition; and since the B-spline basis is local (see equation 2.8), the matrix is banded, i.e.

$$(4.6) \quad g_{ij} \neq 0 \quad \text{if and only if } |i-j| < k, \\ \text{for } 1 \leq i, j \leq m+k.$$

Consequently, only the lower triangle of the band of G need be computed or stored, requiring

$$(4.7) \quad (m+k) \cdot k \text{ locations.}$$

The elements of G and \underline{B} are computed in the following manner:

- 0) Initialize G and \underline{B} to zero. Let $q = 1$.
- 1) Determine an interval of Δ such that

$$(4.8) \quad i \cdot h \leq t_q < (i+1) \cdot h.$$

2) Evaluate those basis functions

$$(4.9) \quad N_{i+1,k}(t_q), \dots, N_{i+k,k}(t_q)$$

which can be nonzero at t_q using a table of their piecewise polynomials (Appendix II).

3) Compute the terms of sums (3.5-6) which depend on the point t_q and are not trivially zero (Figure 4.1). Add them to the appropriate elements of B and of the lower triangle of G :

$$\begin{aligned}
 g_{i+1,i+1} &= g_{i+1,i+1} + N_{i+1,k}(t_q) \cdot N_{i+1,k}(t_q) \\
 g_{i+2,i+1} &= g_{i+2,i+1} + N_{i+2,k}(t_q) \cdot N_{i+1,k}(t_q) \\
 &\vdots \\
 &\vdots \\
 (4.10) \quad g_{i+k,i+k} &= g_{i+k,i+k} + N_{i+k,k}(t_q) \cdot N_{i+k,k}(t_q) \\
 b_{i+1} &= b_{i+1} + y_q \cdot N_{i+1,k}(t_q) \\
 &\vdots \\
 &\vdots \\
 b_{i+k} &= b_{i+k} + y_q \cdot N_{i+k,k}(t_q).
 \end{aligned}$$

4) Let $q = q+1$; if data are exhausted then quit; otherwise go to (1).

	X	X	X	X	X
G:	X	X	X	X	X	B: X
	X	X	X	X	X	X	X
	X	X	X	X	X	X	X	X
i+1	.	X	X	X	+	+	+	+	.	.	.	+
.	.	.	X	X	+	+	+	+	X	.	.	+
.	.	.	.	X	+	+	+	+	X	X	.	+
i+k	+	+	+	+	X	X	X	+
	X	X	X	X	X	X	X
	X	X	X	X	X	X

Figure 4.1: For order $k = 4$, the elements of G and B depending on a given data point satisfying (4.8) are indicated by "+"; trivially zero elements are indicated by "."; and other non-zero elements are indicated by "X".

For an n -dimensional problem, neglecting low order terms, the formation of the normal equations requires

$$(4.11) \sim M \cdot k^2 \text{ multiplications}$$

to evaluate the non-vanishing basis functions,

$$(4.12) \sim \frac{1}{2} \cdot M \cdot k^2 \text{ multiplications}$$

to compute G , and

$$(4.13) \sim n \cdot M \cdot k \text{ multiplications}$$

to compute B . Note that these work estimates are independent of the number of knots and that the work required to compute G is independent of the dimension n .

The normal equations

$$(4.14) G \cdot \underline{A}^* = \underline{B}$$

are solved by a band $L \cdot D \cdot L^T$ factorization. Since G is positive definite and symmetric, there exists a unique factorization of G in the form

$$(4.15) G = L \cdot D \cdot L^T$$

where D is a positive diagonal matrix and L is a lower triangular matrix with unit diagonal [9]. If the relations [13]

$$(4.16) \quad l_{ij} = (g_{ij} - \sum_{q=\max(1,i-k+1)}^{j-1} d_{qq} \cdot l_{iq} \cdot l_{jq}) / d_{jj}$$

for $1 \leq j \leq i \leq m+k$

$$(4.17) \quad d_{ii} = g_{ii} - \sum_{q=\max(1,i-k+1)}^{i-1} d_{qq} \cdot l_{iq}^2$$

for $1 \leq i \leq m+k$

are applied in a proper order, the elements of L and D may be computed from those of G. Using this factorization, the solution to the normal equations is obtained in two steps -- the forward solution for the temporary W

$$(4.18) \quad \tilde{W} = L^{-1} \cdot \tilde{B},$$

and the backward solution for least squares coefficients

$$(4.19) \quad \tilde{A}^* = (L^T)^{-1} \cdot D^{-1} \cdot \tilde{W}.$$

Both steps involve the solution of simple triangular systems.

Band L D L^T factorization has three convenient properties. The diagonal of L is 1; the factor L has the same nonzero structure as the lower triangle of the matrix G; and the element g_{ij} is referenced only once and then it is used to compute l_{ij} (or d_{ii} if $i = j$). Consequently the matrix G may be replaced in storage by the factors L and D as they are computed. To conserve storage only the lower triangles of the bands of L and G are stored. In the code of Appendix I

the two matrices are stored as a vector in row order.
For cubic splines ($k = 4$) the matrix G would be stored as

$$(4.20) \quad g_{11} \quad g_{21}g_{22} \quad g_{31}g_{32}g_{33} \quad g_{41}g_{42}g_{43}g_{44} \quad g_{52}g_{53}g_{54}g_{55}$$

and later replaced in storage by the factorization

$$(4.21) \quad d_{11} \quad l_{21}d_{22} \quad l_{31}l_{32}d_{33} \quad l_{41}l_{42}l_{43}d_{44} \quad l_{52}l_{53}l_{54}d_{55}.$$

This method requires no more storage than that needed for the lower triangle of the band of G and requires far fewer operations than dense Gaussian elimination. The resulting FORTRAN IV subroutine is not significantly larger than a Gaussian elimination routine [9]. For an n -dimensional curve fit, neglecting low-order terms, the factorization of G requires

$$(4.22) \quad \sim \frac{1}{2} \cdot (m+k) \cdot k^2 \text{ multiplications,}$$

the forward solve requires

$$(4.23) \quad \sim n \cdot (m+k) \cdot k \text{ multiplications,}$$

and the backward solve requires

$$(4.24) \quad \sim n \cdot (m+k) \cdot k \text{ multiplications.}$$

Note that the work required to solve the normal equations is independent of the amount of data and that the work required to factor G is independent of the dimension n .

Since the number of data points M is normally much greater than the number of parameters $m+k$, most of the work is in setting up the normal equations (4.11-13) rather than in their solution (4.22-24).

Once the basis coefficients have been computed, the spline may be evaluated for display on the terminal. For computational efficiency the spline should be evaluated from its piecewise polynomial representation, but it was computed in the B-spline representation. Given t and i satisfying

$$(4.25) \quad i \cdot h \leq t < (i+1) \cdot h,$$

the B-spline representation of the spline is given by (equation 2.10)

$$(4.26) \quad S(t) = \sum_{j=i+1}^{i+k} A_j N_{jk}(t)$$

and the piecewise polynomial for the B-spline is given by (Appendix II)

$$(4.27) \quad N_{jk}(t) = \sum_{\ell=0}^{k-1} C_{j-i,\ell}^N \cdot (t - i \cdot h)^\ell \quad \text{for } i+1 \leq j \leq i+k$$

Consequently the coefficients of the piecewise polynomial can be computed from

$$(4.28) \quad C_{i\ell} = \sum_{j=i+1}^{i+k} C_{j-i,\ell}^N \cdot A_j \quad \text{for } 0 \leq i \leq m, \quad 0 \leq \ell \leq k-1.$$

Neglecting lower order terms, the conversion of the entire spline requires

(4.29) $\sim m \cdot k^2$ multiplications.

5. Real Time Calculations

The algorithm of the previous section is neither real-time nor compact. For the sample "Splines" curve and fit of Figure 5.1 (order $k = 4$, dimension $n = 2$, number of data points $M = 2500$, and number of knots $m = 46$), the calculation requires ten seconds of PDP-10 processor time, $3 \cdot 2500$ locations for data storage, $6 \cdot 50$ locations for storage of the least squares arrays G and B , and $4 \cdot 47$ locations for storage of the piecewise polynomial. Clearly, both speed and storage requirements are excessive for a local graphics processor.

With small modifications, the algorithm can be made both compact and real-time. Storage requirements are reduced by eliminating storage of the data and of unneeded elements in the least squares matrices, and real-time response is achieved by performing nearly all of the least squares calculation while the curve is being drawn. In fact the fit for the first part of the curve can be displayed before the entire curve has been drawn, requiring somewhat more computation but far less storage than for a sequential implementation.

The calculation divides naturally into three major processes -- the formation of the normal equations, the factorization and forward solution of the normal equations,

The image shows the word "Splines" written in a cursive script. The letters are composed of discrete points connected by thin lines, illustrating the concept of splines. The word is oriented vertically on the page.

Figure 5.1a: "Splines" data 2500 discrete points taken from a ten bit resolution graphics tablet.

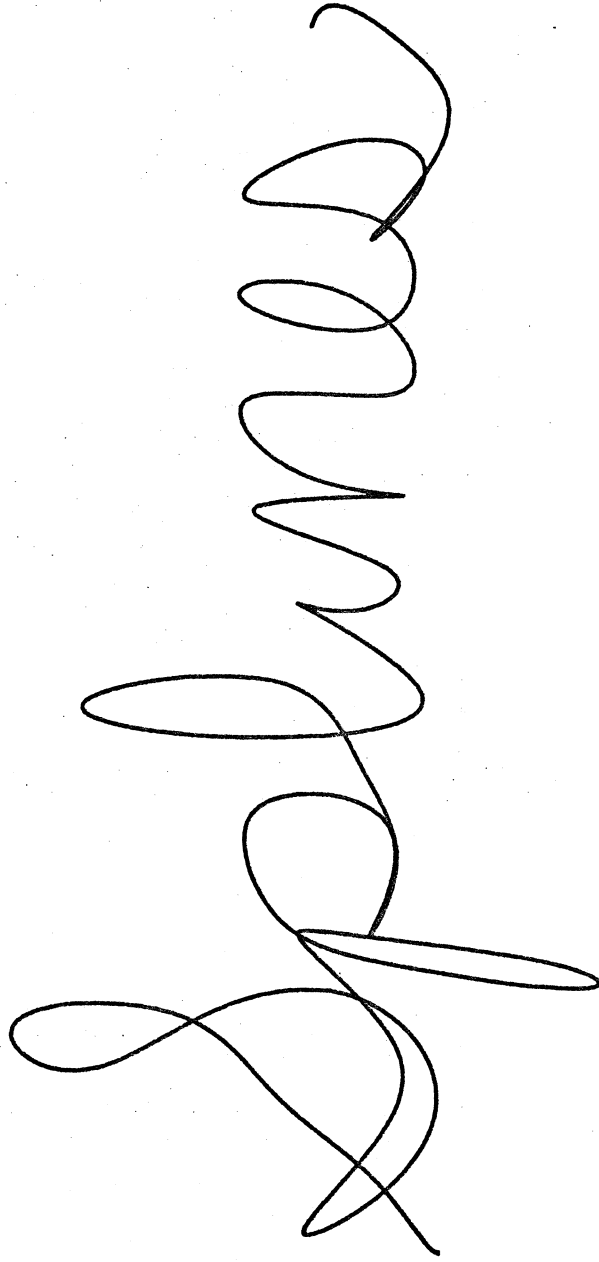


Figure 5.1b: Cubic spline fit to "Splines" data, $m = 46$.

and the back solution for the basis coefficients and their conversion to a piecewise polynomial. These processes need not be performed sequentially. In this section they are pipe-lined so that, at any given point in time, each process will have been completed to the greatest extent possible. In this way the idle time during curve drawing is employed to display the fit as soon as possible.

Of the three processes, the formation of the normal equations requires the most processor time (Table 5.1).

1) Normal equations	$\sim M \cdot k \cdot (3k/2 + n)$ multiplications ~ 80000
2) Factorization and forward solve	$\sim (m+k) \cdot k \cdot (k/2 + n)$ multiplications ~ 800
3) Backsolve and conversion	$\sim n \cdot (m+2) \cdot k^2$ multiplications ~ 1536

Table 5.1: Work estimates for the three major spline least squares processes with numerical values computed for the "Splines" curve.

The computation of the normal equations need not be postponed until the last data point has been acquired. As each data point is received, the nonzero terms of the sums

(3.5-6) to which that point contributes may be computed and added to the appropriate elements of the least squares arrays. Unless the data point is needed for display or for later comparison with the fit, it need not be stored at all. A data point satisfying

$$(5.1) \quad i \cdot h \leq t < (i+1) \cdot h$$

affects only rows $i+1$ to $i+k$ of B and of the lower triangle of G (see Figure 4.1); consequently, after the last point of interval i has been acquired, the values of rows 1 to $i+1$ are final and are available for further processing.

As each row of G is determined, its factorization may be computed from (4.16-17). The elements l_{ij} and d_{ii} of the factorization are computed in left-to-right order, and for efficient storage utilization they replace the corresponding elements of G . After the factorization is complete to row i , the forward solve may be completed up to element w_i and the elements of W may replace the corresponding elements of B (see Figure 5.2).

The third process -- backsolve and conversion -- cannot be begun until both the forward solution and the factorization are complete. Neither the forward solution nor the factorization can be completed until all of the data have been received; therefore the algorithm as stated will not allow computation of the value of any basis

<p>G: d</p> <p> ℓ d</p> <p> ℓ ℓ d</p> <p> ℓ ℓ ℓ d</p> <p> ℓ ℓ ℓ d</p> <p> ℓ ℓ ℓ d</p> <p>i+1 g g g g</p> <p> . g g g g</p> <p> . g g g g</p> <p>i+k g g g g</p>	<p>B: w</p> <p> w</p> <p> w</p> <p> w</p> <p> w</p> <p> w</p> <p> w</p> <p> b</p> <p> b</p> <p> b</p> <p> b</p>
---	---

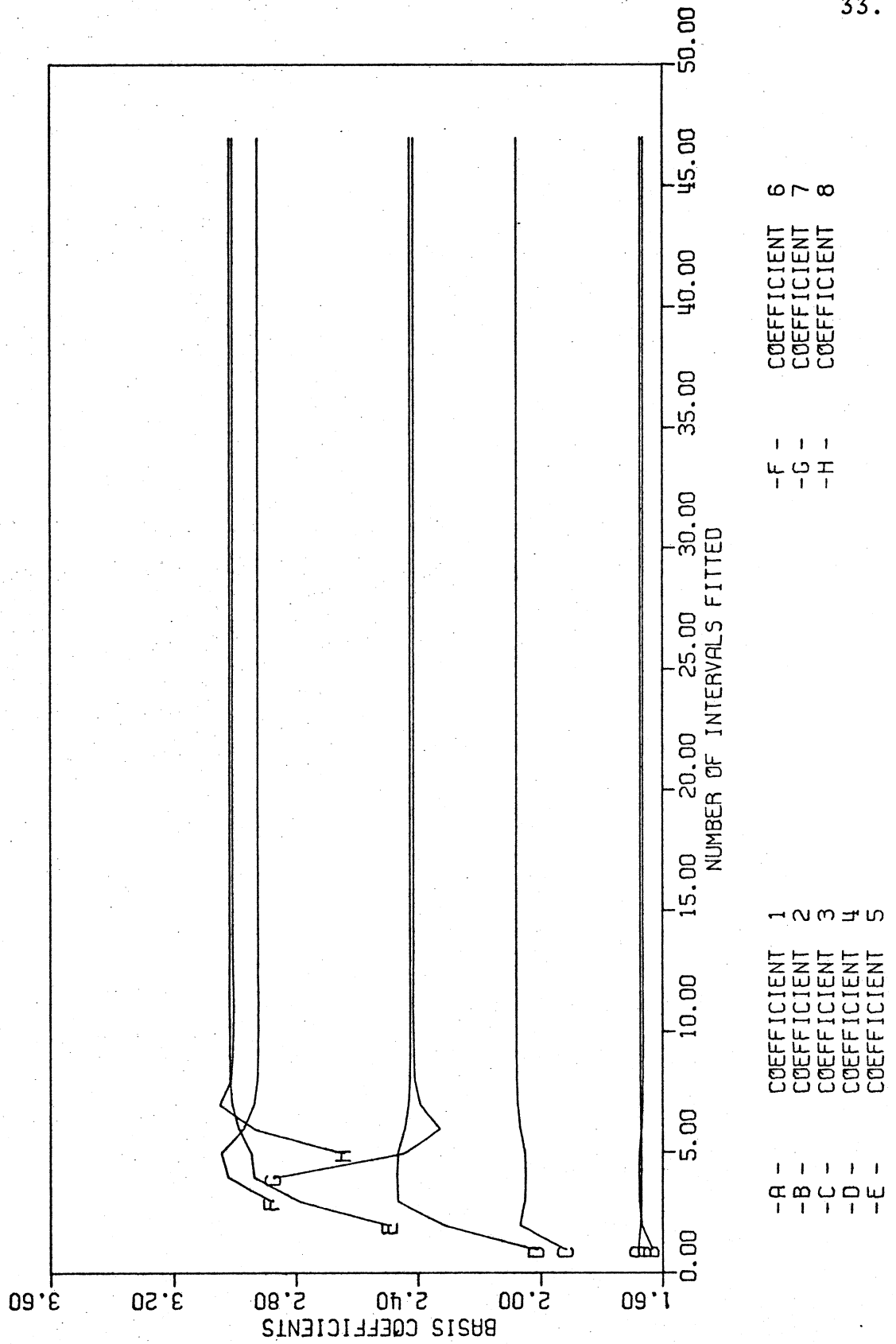
Figure 5.2: The least squares matrices and their factorization during processing of data point t_q satisfying equation 4.1.

coefficient before all of the data have been acquired. Furthermore, since the value of each point on a least squares spline curve depends on the values of all of the data, no method exists for computing any part of a least squares curve without knowing the entire set of data.

However, as the experiment of Figure 5.3 suggests, one may compute approximate values for some of the basis coefficients without the complete set of data. In Figure 5.3, for the "Splines" fit of Figure 5.1, the values of the first eight B-spline basis coefficients were computed using data in the first I intervals and the computed values are plotted for values of I from 1 to 47. As the fraction of data fitted increases, the values of the coefficients approach those computed from the full set of data. For these first eight coefficients, three digit accuracy may be obtained using data in the beginning ten to fifteen intervals only.

This effect is a result of the local convergence property of splines. Stated informally, local convergence implies that the value of a spline at a given point depends mostly on the data in the neighborhood of that point, and that the value of a basis coefficient depends mostly on the data in the neighborhood of the intervals where the corresponding basis function is non-zero. Data far from a given point on a curve affect only slightly the value of

Figure 5.3: Convergence of least squares coefficients to final values as fraction of data fitted increases. Data are from "Splines" curve of Figure 5.1a.



the least squares fit near that point. More precise statements of this property and proofs for some special cases may be found in Powell [14] and Strang and Fix [20]. A proof for the general case appears in Lewis [11].

In interactive graphics applications, where two or three digit accuracy is quite adequate, local convergence allows computation of the spline fit to the first part of curve data before the entire set of data is acquired. Given the first $i+1$ rows of the forward solve W and of the factorization $L\backslash D$, the backsolve for rows 1 to $i+1$ may be performed. For any standard of accuracy, for i sufficiently large, and for some choice of a positive integer $\text{lag} < i+1$, coefficients

$$(5.2) \quad A_{i-\text{lag}+1}, \dots, A_i, A_{i+1}$$

will not be sufficiently accurate, but the coefficients

$$(5.3) \quad A_1, \dots, A_{i-\text{lag}}$$

will be sufficiently accurate and may be converted to a piecewise polynomial for display (Figure 5.4). On subsequent backsolves the coefficients of (5.3) need not be computed again.

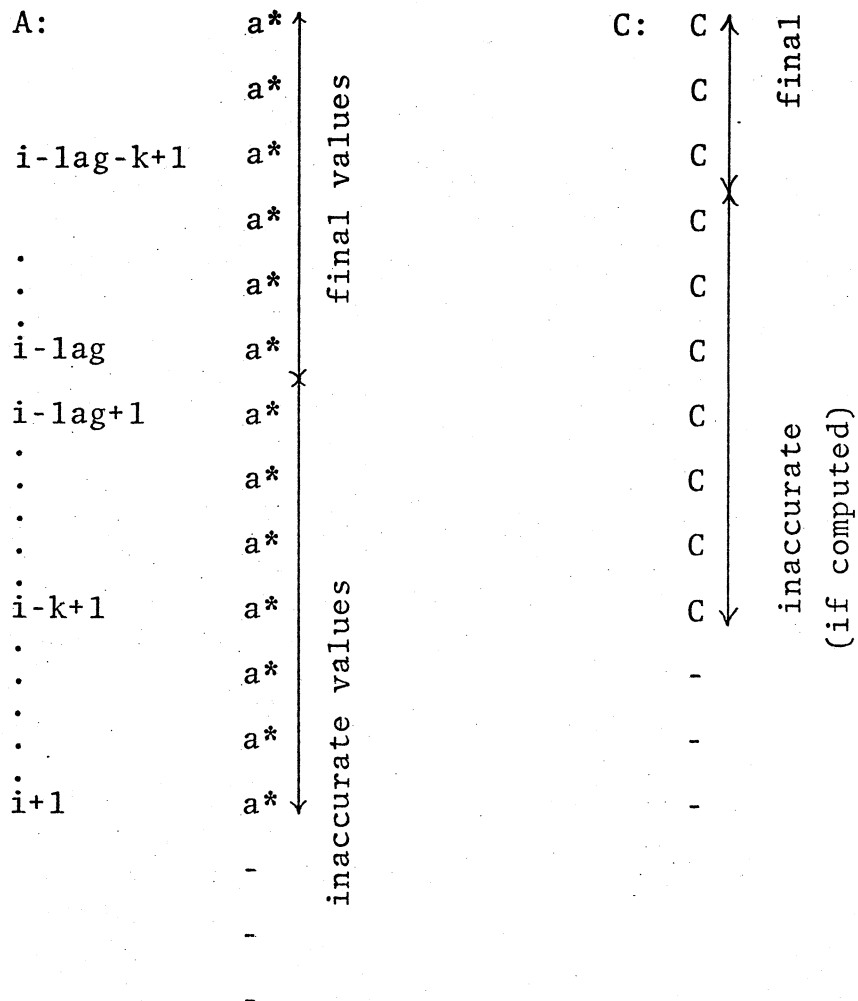


Figure 5.4: Backsolve and conversion.

Assuming that a backsolve is performed at least once per interval and neglecting low order terms, each backsolve requires

(5.4) $\text{lag} \cdot k$ multiplications.

Clearly the value of lag should be given the smallest value consistent with working accuracy. For curve data the choice $\text{lag} = 2 \cdot k$ appears to produce reasonable results without requiring excessive computation.

Other than the additional computation in the backsolve, the computation required for this algorithm is the same as that for the sequential algorithm of Section 4. Assuming that the backsolve is performed once per interval, the computational penalty for obtaining the fit before completion of the curve lies entirely in the backsolve and amounts to approximately

(5.5) $[m \cdot \text{lag} - (m+k)] \cdot k$ multiplications

over the sequential algorithm.

If the backsolve is performed once per interval, this scheme requires less storage than the sequential algorithm. After the value for $A_{i-\text{lag}}$ has been computed, rows

(5.6) $1, 2, \dots, i-\text{lag}$

of the least squares matrices are no longer required and

their storage may be freed for other uses. Only

(5.7) $lag \cdot (k+1)$ locations

of the matrices G and B need be kept. Therefore, unless the basis coefficients or the piecewise polynomial must be saved, the storage required for the algorithm is independent of the length of the curve.

6. Closed Curves

For a general curve-design system, both open and closed curve capabilities are required. Objects such as bowling pins, television tubes, and bottles contain closed cross-section curves, and the designer must be able to create and manipulate these closed curves as easily as open curves. Furthermore, the system to fit and manipulate these closed curves should meet the same goals as the open curve system -- compactness and real-time response.

The entire development of Sections 3 to 5 generalizes to closed curves. The closed spline curve is simply an open curve whose endpoints join smoothly, i.e. an open curve is also a closed curve if it satisfies the periodic end conditions

$$(6.1) \quad D^{\ell}S(0^+) = D^{\ell}S(a^-) \quad \text{for } 0 \leq \ell \leq k-2.$$

For $m > k$ this derivative constraint is equivalent to the constraint

$$(6.2) \quad \tilde{A}_i = \tilde{A}_{i+m+1} \quad \text{for } 1 \leq i \leq k-1$$

on the B-spline basis coefficients.

Except for the additional constraint (6.2), the calculation of the periodic least squares spline curve follows the development of Section 3. Since the curve fitting problem reduces to a number of one variable least

squares problems, only the one variable case will be treated here.

Given the order k , the knots

$$(6.3) \quad \Delta: 0, h, \dots, m \cdot h, (m+1) \cdot h = a$$

and the data

$$(6.4) \quad X = \{(t_q, x_q): 1 \leq q \leq M\},$$

the closed (or periodic) least squares spline is that spline $S_k^\Delta(t)$ which is closest to the data in that it minimizes the distance functional

$$(6.5) \quad d(S, X) = \sum_{q=1}^M [S(t_q) - x_q]^2$$

over all splines of given order and knots which satisfy the smoothness constraint (6.2). As in Section 3 the distance functional $d(S, X)$ may be written as a quadratic function of \tilde{A}

$$(6.6) \quad F(\tilde{A}) = \tilde{A}^T \cdot G \cdot \tilde{A} - 2 \cdot \tilde{A}^T \cdot B + C$$

where G , B , and C are defined in Section 3.

The constraint (6.2) may be eliminated by substitution of A_{i-m-1} for A_i if $i > (m+1)$ in equation (6.6). After this substitution (6.6) becomes

$$(6.7) \quad F(\tilde{A}) = \tilde{A}^{\circ T} \cdot \tilde{G} \cdot \tilde{A} - 2 \cdot \tilde{A}^{\circ T} \cdot \tilde{B} + C$$

where the periodic Gram matrix \tilde{G}

$$(6.8) \quad \tilde{G} = [g_{ij}] = \left[\begin{array}{ll} g_{ij} + g_{i+m+1, j+m+1} & \text{for } 1 \leq j, i \leq k-1 \\ g_{j+m+1, i} & \text{for } 1 \leq j \leq k-1 \\ & \text{and } m-k+3 \leq i \leq m+1 \\ g_{j, i+m+1} & \text{for } 1 \leq i \leq k-1 \\ & \text{and } m-k+3 \leq j \leq m+1 \\ g_{i, j} & \text{otherwise} \end{array} \right]$$

and the periodic vector \tilde{B}

$$(6.9) \quad \tilde{B} = [b_i] = \left[\begin{array}{ll} b_i + b_{i+m+1} & \text{for } 1 \leq i \leq k-1 \\ b_i & \text{otherwise} \end{array} \right]$$

are obtained by "folding" the non-periodic matrices G and B (Figure 6.1). This "folding" computation requires no multiplications and

$$(6.10) \quad \sim \frac{1}{2} k^2 \text{ additions.}$$

The coefficients

$$(6.11) \quad A_i = \left\{ \begin{array}{ll} A_i & \text{for } 1 \leq i \leq m+1 \\ A_{i-m-1} & \text{for } m+2 \leq i \leq m+k \end{array} \right\}$$

of the corresponding open curve are the periodic extension

	X	X	X		X
G:	X	X	X	X	B:	X
	X	X	X	X	X		X
	.	X	X	X	X	X		X
	.	.	X	X	X	X	X	.	.	.		X
	.	.	.	X	X	X	X	X	.	.		X
m-k+3	X	X	X	X	1	.		X
m+1	X	X	X	2	3		X
m+2	1	2	4	5		1
m+k	3	5	6		2

Figure 6.1a: The non-periodic matrix G and vector B for $k = 3$, $m = 7$. Nonzero elements are indicated by "X" and zero elements by ".". Numbered elements are the nonzero elements that are added to the non-periodic matrix G to produce the periodic matrix \tilde{G} .

1	4	5	X	.	.	.	1	2		1
$\overset{\circ}{G}: k-1$	5	6	X	X	.	.	.	3		$\overset{\circ}{B}: 2$
	X	X	X	X	X	.	.	.		X
	.	X	X	X	X	X	.	.		X
	.	.	X	X	X	X	X	.		X
	.	.	.	X	X	X	X	X		X
$m-k+3$	1	.	.	.	X	X	X	X		X
$m+1$	2	3	.	.	.	X	X	X		X

Figure 6.1b: The periodic matrix $\overset{\circ}{G}$ and vector $\overset{\circ}{B}$ for $k = 3$, $m = 7$. Nonzero elements from the non-periodic matrices are indicated by "X" and zero elements by ".". The numbered elements are the sums of the corresponding elements in the non-periodic matrix and the elements of the same number in Figure 6.1a.

of the closed curve coefficients \tilde{A}^* .

The periodic Gram matrix $\overset{\circ}{G}$ is well conditioned, positive definite, and symmetric; but it is not banded (Figure 6.1b). While band methods are not efficient for such a matrix, variable bandwidth or envelope methods are quite effective [7]. For the symmetric matrix $\overset{\circ}{G}$, the lower half of the envelope of $\overset{\circ}{G}$ consists of the elements

$$(6.12) \{g_{ij}^{\circ} : f_i \leq j \leq i, 1 \leq i \leq m+1\} \quad \text{where}$$

$$(6.13) f_i \equiv \min j \text{ such that } g_{ij}^{\circ} \neq 0$$

Only the envelope of the matrix is ever stored or computed. The elements of the lower triangle of the envelope of G are stored in a real vector V^G and accessed through the integer vector U^G

$$(6.14) g_{ij}^{\circ} = V_{j+u_i}^G$$

with

$$(6.15) u_i^G = \sum_{j=1}^i (j - f_j).$$

Storage of the periodic matrix G requires fewer than

$$(6.16) \quad (m+1) \cdot (2k-1) \text{ locations for } V^G \\ \text{and } (m+1) \text{ locations for the pointers } U^G.$$

exists. From the relations [7]

$$(6.18) \quad \ell_{ij} = (g_{ij} - \sum_{q=\max(f_i, f_j)}^{i-1} d_{qq} \cdot \ell_{iq} \cdot \ell_{jq}) / d_{jj}$$

for $1 \leq j \leq i \leq m+k$

and

$$(6.19) \quad d_{ii} = g_{ii} - \sum_{q=f_i}^{i-1} d_{qq} \cdot \ell_{iq}^2 \quad \text{for } 1 \leq i \leq m+k$$

the elements of D and L may be computed. Just as the factorization of G for the non-periodic case has the same bandwidth as G, so the factorization of $\overset{\circ}{G}$ has the same envelope as $\overset{\circ}{G}$. Consequently, the factorization of $\overset{\circ}{G}$ may replace $\overset{\circ}{G}$ in memory just as before. After a forward solve

$$(6.20) \quad \overset{\circ}{\tilde{W}} = L^{-1} \cdot \overset{\circ}{B}$$

for $\overset{\circ}{\tilde{W}}$, the solution $\overset{\circ}{A}^*$ is obtained from the back solution

$$(6.21) \quad \overset{\circ}{\tilde{A}}^* = (L^T)^{-1} \cdot D^{-1} \cdot \overset{\circ}{\tilde{W}}.$$

For an n-dimensional problem and neglecting low order terms, the factorization requires

$$(6.22) \quad 2 \cdot (m+1) \cdot k^2 \text{ multiplications,}$$

the forward solve requires

(6.23) $2 \cdot n \cdot (m+1) \cdot k$ multiplications,

and the back solve requires

(6.24) $2 \cdot n \cdot (m+1) \cdot k$ multiplications.

In practice the curve would be drawn and fit as an open curve (Figure 6.3) and the ends would be joined later. If the designer were so kind as to draw a curve whose length was an integer multiple of the knot spacing h , i.e.

(6.25) $a = h \cdot (m+1)$

then the fit could be obtained easily by "folding" the non-periodic matrices and solving the resulting linear system (Figure 6.4). Unfortunately, if the designer were to draw a curve not satisfying equation (6.25), kinks and bulges could be expected if the closed curve were calculated from the open curve fit (Figure 6.5) by folding (Figure 6.6).

If equation (6.25) is not satisfied, to obtain an exact periodic least squares fit the entire computation could be repeated with

(6.26) $h^{\text{new}} = \frac{a}{m+1}$.

This fit would require considerable computation -- more than the original open curve fit itself -- and might introduce unacceptable delay in an interactive system. Furthermore,

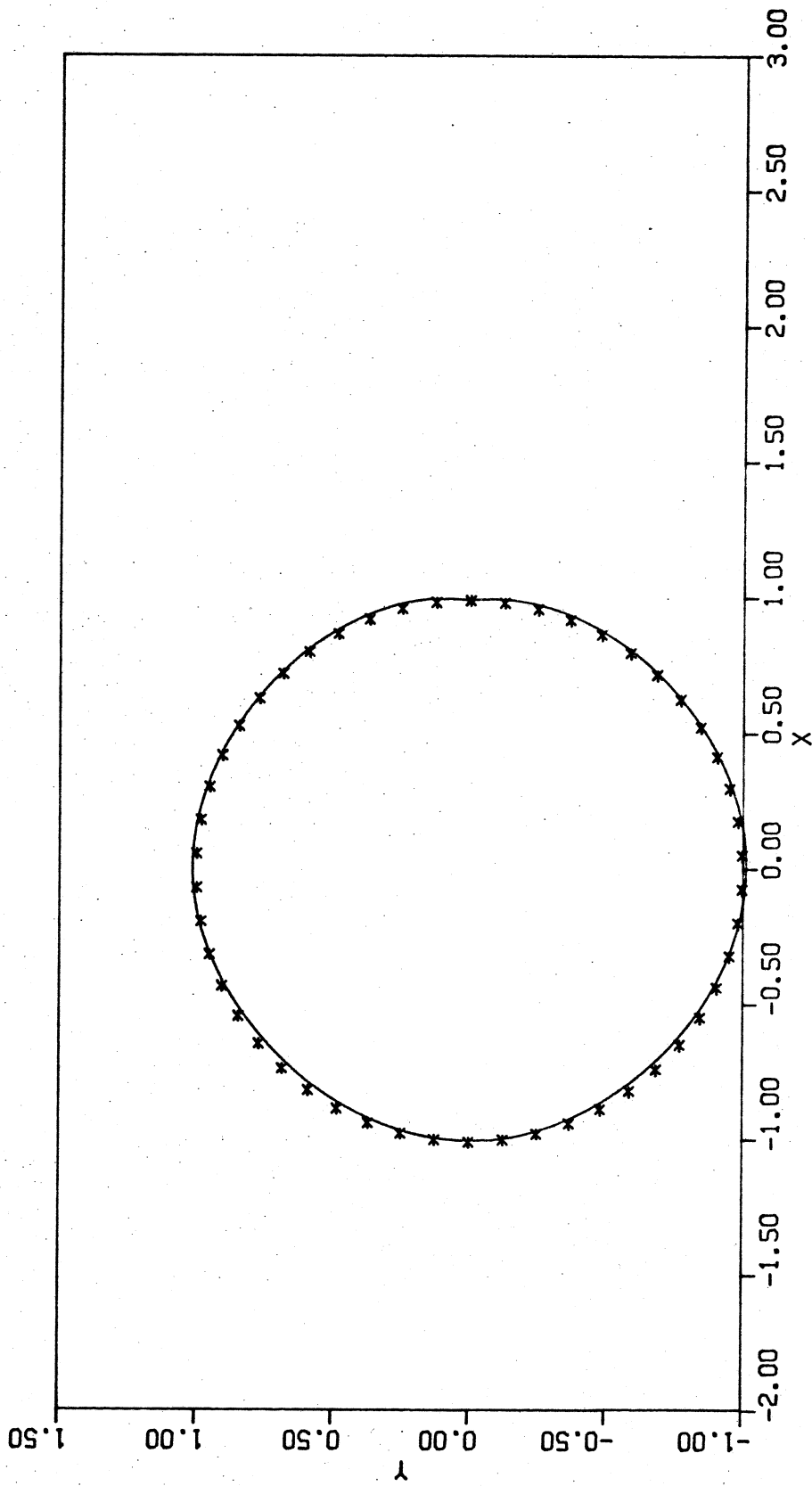


Figure 6.3: Open curve fit, $5 \cdot h = a$, $m = 4$, $k = 4$.

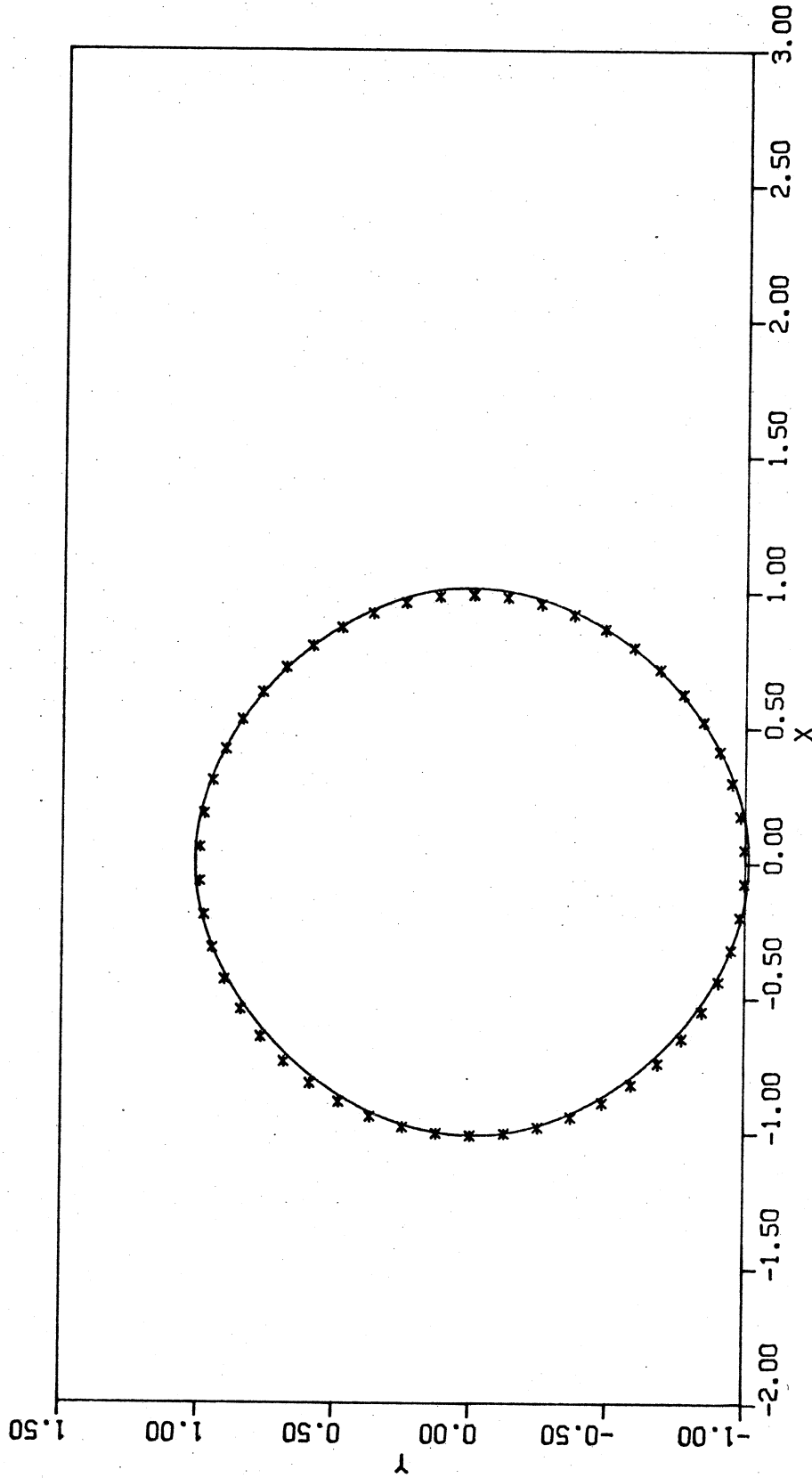


Figure 6.4: Closed curve fit, $5 \cdot h = a$, $m = 4$, $k = 4$.

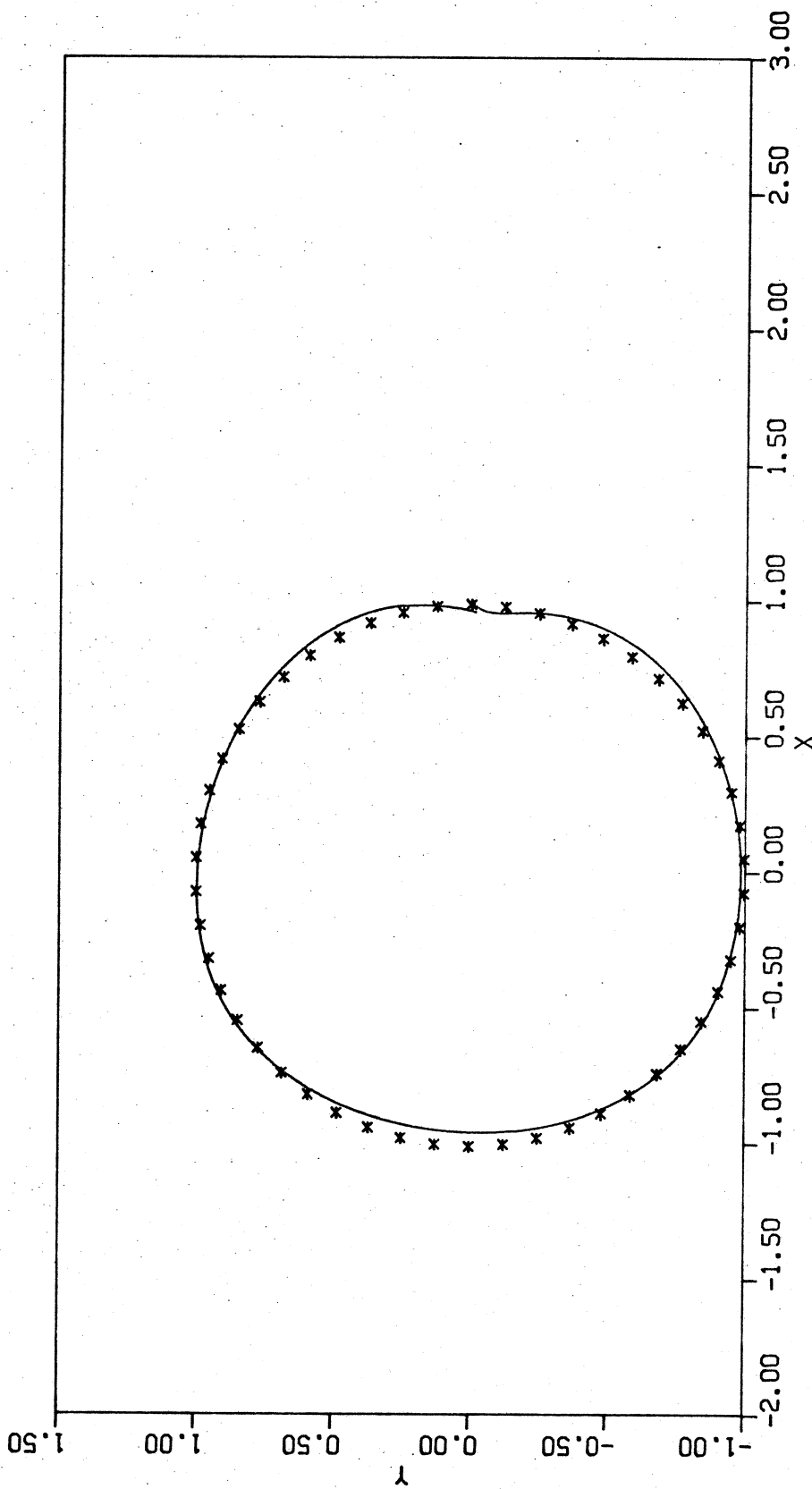


Figure 6.5: Open curve fit, $4.1 \cdot h = a$, $m = 4$, $k = 4$.

6.5

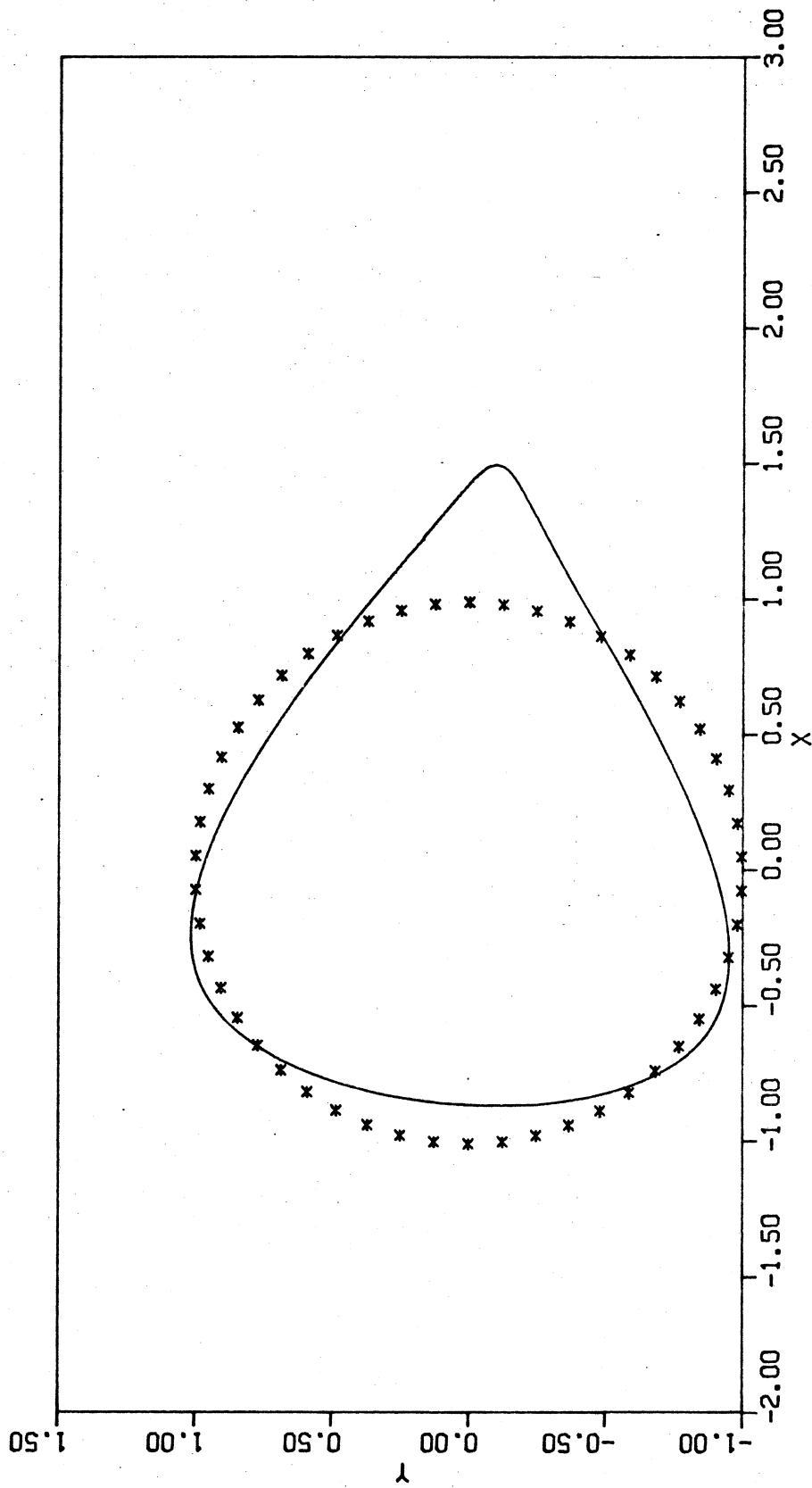


Figure 6.6: Closed curve fit, $4.1 \cdot h = a$, $m = 4$, $k = 4$.

all of the data would have to be stored. To avoid much of this delay and the requirement of data storage, an approximate closed curve fit can be obtained by fitting a closed curve to points evaluated from the open curve fit (Figure 6.5 illustrates the open curve fit and Figure 6.7 is a closed curve fit to data from that fit). If the evaluation points are chosen carefully, such as at the Gaussian quadrature points, relatively few evaluations are necessary and the time required for a closed curve fit is a small fraction of the computation time for the exact fitting procedure.

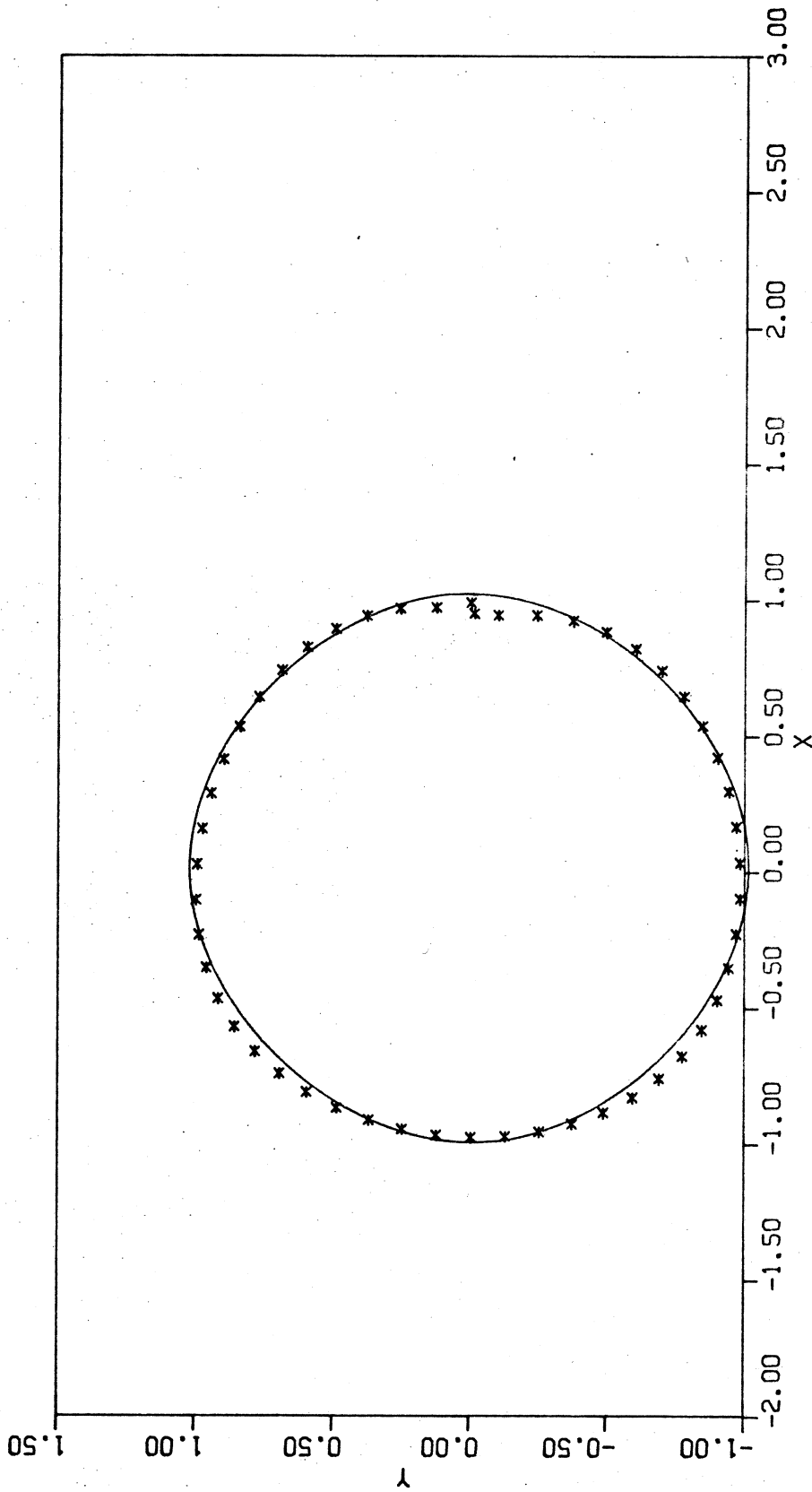


Figure 6.7: Closed curve fit using data obtained from fit of Figure 6.5,
 $5 \cdot h = a$, $m = 4$, $k = 4$.

7. Acknowledgments

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Appendix I: FORTRAN IV Codes

The algorithm of this paper has been implemented in FORTRAN IV and tested with the PDP-10 F4 compiler, the OS370 G compiler, and the CDC RUN compiler. The codes are close to the ASA standard for FORTRAN but do violate the standards for integer expressions in array indices. There are two basic packages -- CURVS, the curve fitting package, and SPLSQ, a package for weighted least squares splines in one dimension. Both packages use the modules in PSQARS (primitive least squares operations) and ENSOLV (incremental envelope linear system solution).

All of the code is well-documented and highly modular. The code was designed more to illustrate the implementation of the algorithm than for efficiency. For maximum efficiency, the subroutine parameters should be placed in COMMON blocks and some of the subroutine code should be placed in-line.

Use of the packages is illustrated by a set of drivers. All drivers include sample data files and part of the resulting output so that the proper functioning of the codes may be verified conveniently. There are three drivers for the CURVS package, each illustrating a different mode of use. These three modes are explained at the beginning of the CURVS module. There is one simple data fitting driver for the SPLSQ subroutine.

The SPLSQ subroutine is well-protected against user error in that it checks subroutine parameters and in case of error returns an explanatory flag; but because of its structure, the CURVS package could not be so well-protected. Variables and arrays passed from subroutine to subroutine (for example, JG, G, LOW, LOWF) should not be modified, and the subroutines must be called in the specified order.

Machine-readable copies of the codes are available
from

Numerical Software
Department of Computer Science
Yale University
New Haven, Connecticut 06520.

```

C *****
C * Module SPLSQ
C * Least Squares Polynomial Splines
C * User Subroutine
C * 22 Sep 74
C *****
C Entries: SPLSQ
C External: from module PSOARS; SETUPS,ADDON,FOLD,CONVRT
C          from module ENSOLV; FACTOR,SOLVE
C
C Reference: S.C. Eisenstat, J.w. Lewis, M.H. Schultz,
C "A Real-Time Algorithm for Least Squares Splines and
C Its Application in Computer Aided Geometric Design",
C research report #29, Department of Computer Science,
C Yale University
C
C Subroutine SPLSQ fits a periodic or non-periodic least squares
C spline to weighted data XD(ND), WD(ND), Y(LY,NY). The least squares
C spline is returned as the basis coefficients A(LA,NY) and as the
C piecewise polynomial CPP(LCK,ICM,NY). The subroutine is well
C protected against user error. If parameters are incorrect, SPLSQ
C returns with an explanatory error flag.
C *****Subroutines and Calling Sequences*****
C
C 1) Unless otherwise indicated, variable types are
C    IMPLICIT REAL(A-H, O-Z)
C    IMPLICIT INTEGER(I-N)
C
C 2) Value variables (V) pass a value to the subroutine and must
C    have been set by the user or by a previously called subroutine.
C    Result variables (R) return results to the calling subprogram.
C
C    CALL SPLSQ( JPER, K, MK,AA,BB, ND,XD,WD, NY,LY,YD,
C              1 JG,IG,G,NG, LA,A, LCK,ICM,CPP,HKNOT, LFLAG)
C
C V JPER - Nonzero to indicate periodic problem
C V K - Degree +1 of spline
C V MK - Number of knots (including endpoints)
C V AA - Left endpoint of interval
C V BB - Right endpoint of interval
C
C V ND - Number of evaluation points
C V XD(ND) - Evaluation points
C V WD(ND) - Weight for each evaluation point
C          If WD(I)<0 then points are equally weighted
C V NY - Number of sets of ordinate values
C V LY - Dimension of YD
C V YD(LY,NY) - NY sets of ND ordinate values
C
C V JG(LA) - Index array for grammian
C V IG - Dimension of G
C V LG - Storage for grammian (profile only is stored)
C V G(LG) - G(I,J) = G( JG(I) + J ) for I le J
C
C V NG - Number of locations used in G
C
SPLS 1
SPLS 2
SPLS 3
SPLS 4
SPLS 5
SPLS 6
SPLS 7
SPLS 8
SPLS 9
SPLS 10
SPLS 11
SPLS 12
SPLS 13
SPLS 14
SPLS 15
SPLS 16
SPLS 17
SPLS 18
SPLS 19
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SPLS 50
SPLS 51
SPLS 52
SPLS 53
SPLS 54
SPLS 55

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```

C V LA - Dimension of A
C R A(LA,NY) - Basis coefficients
C V LCK,ICM - Dimensions of CPP
C R CPP(LCK,ICM,NY) - Piecewise polynomial
C R HKNOT - Knot spacing (for use by EVAL)
C
C R LFLAG - Status (either -1, 0 or a sum of the positive codes)
C
C -1 - Problem is ill posed - badly distributed data
C 0 - Successful fit
C 1 - K LT 1
C 2 - K GT KMAX or K GT LCK
C 4 - MK LT 2
C 8 - (MK-1) GT ICM
C 16 - MK LE K and JPER ne 0
C 32 - AA LE BB
C 64 - NY LT 1
C 128 - NA GT LA
C 256 - NG GT LG
C 512 - N GT ND insufficient data
C *****
C SUBROUTINE SPLSQ( JPER, K, MK,AA,BB, ND,XD,WD, NY,LY,YD,
C DIMENSION XD(NY), WD(NY), YD(LY,NY)
C DIMENSION JG(LA),G(LG), A(LA,NY), CPP(LCK,ICM,NY)
C DIMENSION CB(36)
C DATA KMAX/6/
C
C Compute storage used and check parameters
C
C NA = MK-K-2
C N = NA
C IF(JPER.NE.0) N = MK-1
C NG = NA*K - K*(K-1)/2
C
C Error checks
C LFLAG = 0
C IF( K.LE.0 )
C IF((K.GT.KMAX).OR.(K.GT.LCK) )
C IF( MK.LT.2 )
C IF( (MK-1).GT.ICM )
C Additional restriction for periodic problems
C IF(JPER.EQ.0) GO TO 20
C IF( MK.LE.K )
C K1 = K-1
C DO 10 I=1,KM1
C NG = NG + MAX0( 0, I-JF -K )
C IF( AA .GE. BB )
C IF( NY .LE. 0 )
C IF( NA .GT. LA )
C IF( NG .GT. LG )
C IF( N .GT. ND )
C IF( LFLAG.NE.0 ) RETURN
C
C 10
C 20
C
C Setups
C
SPLS 56
SPLS 57
SPLS 58
SPLS 59
SPLS 60
SPLS 61
SPLS 62
SPLS 63
SPLS 64
SPLS 65
SPLS 66
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SPLS 90
SPLS 91
SPLS 92
SPLS 93
SPLS 94
SPLS 95
SPLS 96
SPLS 97
SPLS 98
SPLS 99
SPLS 100
SPLS 101
SPLS 102
SPLS 103
SPLS 104
SPLS 105
SPLS 106
SPLS 107
SPLS 108
SPLS 109
SPLS 110

```

```

C *****
C * SPLSQ driver
C * Least Squares Polynomial Splines
C * 24 Aug 74
C *****
C Externals: from module SPLSQ
C from module PSCARS: EVAL
C
C Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
C "A Real-Time Algorithm for Least Squares Splines and
C Its Application in Computer Aided Geometric Design",
C research report #29, Department of Computer Science,
C Yale University
C
C This program is a driver for the SPLSQ spline fitting subroutine.
C It reads data from a file and prints the piecewise polynomial
C for the fit.
C
C For LFLAG error messages, see module SPLSQ
C
C Sample data file (sine and cosine)
C... 0.000000 6.283185 JPER,K,MK,ND,NY (211,12,12,11)
C... 0.000000 0.000000 Interval of fit (2F10.6)
C... 0.628319 0.587785 Data XD, YD(1), ...,XD(NY)
C... 1.256637 0.951057 (8F10.6)
C... 1.849956 0.951057 JPER - end conditions
C... 2.513274 0.951057 periodic for JCON NE 0
C... 3.141593 0.000000 -1.000000 K - Degree + 1
C... 3.769911 -0.587785 -0.309017 ND - Number of data points XD
C... 4.398230 -0.951057 -0.309017 NY - Number of YD values
C... 5.026548 -0.951057 0.309017 at each point XD
C... 5.654867 -0.587785 0.609017
C... 6.283185 -0.000000 1.000000
C
C Results:
C
C SPLINE FIT
C JPER K MK ND NY AA BB
C 0 4 5 11 2 0.0000 6.2832
C
C PIECEWISE POLYNOMIAL... 1
C 0.00124 0.98717 -0.02238 -0.12471
C 1.01332 -0.00626 -0.61005 0.12946
C 0.00000 -0.96452 0.00000 0.12946
C -1.01332 -0.00626 0.61005 -0.12471
C
C PIECEWISE POLYNOMIAL... 2
C 0.99916 0.08489 -0.70341 0.15461
C -0.00385 -0.98048 0.02517 0.12177
C -1.00990 -0.00000 0.59902 -0.12177
C -0.00385 0.98048 0.02517 -0.15461
C
C DIMENSION XD(101), WD(1), YD(120,6)
C DATA NDMAX/101/, LY/120/, NYMAX/6/
C DIMENSION JG(20), A(20,6), G(300), CPP(6,20,6)
C DATA LA/20/, LG/300/, LCK/6/, LCM/20/

```

```

SODR 1
SODR 2
SODR 3
SODR 4
SODR 5
SODR 6
SODR 7
SODR 8
SODR 9
SODR 10
SODR 11
SODR 12
SODR 13
SODR 14
SODR 15
SODR 16
SODR 17
SODR 18
SODR 19
SODR 20
SODR 21
SODR 22
SODR 23
SODR 24
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SODR 35
SODR 36
SODR 37
SODR 38
SODR 39
SODR 40
SODR 41
SODR 42
SODR 43
SODR 44
SODR 45
SODR 46
SODR 47
SODR 48
SODR 49
SODR 50
SODR 51
SODR 52
SODR 53
SODR 54
SODR 55

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```

C Uniform weighting of data
C WD(1) = -1.
C
C Read Data
C
C 100 READ(5,11010) JPER,K, MK, ND, NY, AA, BB
C 11010 FORMAT(2I1,12,13,11/ZF10.4)
C 110200 WRITE(6,11020) JPER,K,MK,ND,NY,AA,BB
C 110200 FORMAT(/,/, Spline Fit /
C 1 JPER K MK ND NY AA BB /
C 2 514,ZF10.4)
C
C 11030 WRITE(6,11030) .AND. (NY.LE.NYMAX) ) GO TO 110
C 11030 FORMAT( ' Bad input parameters ' )
C STOP
C DO 120 I=1,ND
C 120 READ(5,11040) XD(I), (YD(I,J), J=1,NY)
C 11040 FORMAT(8F10.6)
C
C Fit data
C
C 200 CALL SPLSQ( JPER, K, MK,AA,BB, ND,XD,WD, NY,IY,YD,
C 1 JG,LG,G,NG, LA,A, LCK,LCM,CPP,HMESH, LFLAG )
C Check for error return
C IF (LFLAG.EQ.0) GO TO 210
C WRITE(6,12010) LFLAG
C 12010 FORMAT( ' ERROR ',I10)
C STOP
C 210 CONTINUE
C
C Print Piecewise Polynomials
C
C 300 MKM1 = MK-1
C DO 320 I=1,NY
C WRITE(6,13010) I
C 13010 FORMAT( /, ' Piecewise polynomial...',I2)
C DO 310 M=1,MKM1
C WRITE(6,13020) (CPP(J,M,I), J=1,K)
C 13020 FORMAT(6F13.5)
C CONTINUE
C STOP
C END

```

```

1 1 *****
2 2 CURVS driver
3 3 Incremental Least Squares Splines
4 4 Mode 2 Test
5 5 29 Aug 74
6 6 *****
7 7 Externals: from module CURV1,CURV2,CURV3
8 8 from module PSQARS: EVAL
9 9
10 10 Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
11 "A Real-Time Algorithm for Least Squares Splines and
12 Its Application in Computer Aided Geometric Design",
13 research report #29, Department of Computer Science,
14 Yale University
15
16 C This program is a driver for the CURVS incremental curve
17 fitting package. It reads data from a file and prints the
18 basis coefficients and piecewise polynomial for the fit.
19
20 C For LFLAG error messages, see module CURVS
21
22 C A sample data file:
23 K,ND,NE,HKNOT (12,215,F10.3)
24 X,Y (2F10.6)
25
26 C... 1. 9.
27 C... 2. 8.5
28 C... 3. 8.
29 C... 4. 6.5
30 C... 4.5 5.
31 C... 5. 3.
32 C... 2.5
33 C... 3.
34 C... 4.
35 C... 2.
36 C... 1.
37
38 C Resulting Piecewise polynomials (to check operation of program)
39
40 C 1.071615 0.733437 0.051666 -0.011503 For X coordinate
41 4.862930 -0.052865 -0.172635 0.012959
42
43 C 8.858785 0.167781 -0.279348 0.021057 For Y coordinate
44 3.929716 -0.794753 0.131266 0.004390
45
46 DIMENSION Y(2), YO(2)
47 DATA NY/2/
48 DIMENSION A(80,2), B(80,2), JG(80), G(6,80)
49 DIMENSION CPP(6,80,2), CB(6,6)
50 DATA LCK/6/, LA/80/, LCH/80/
51
52 C Fit with Spline
53
54 C Read K,ND,HKNOT and form matrices
55 READ(5,10010) K, ND,NE, HKNOT
56 WRITE(6,10020) K, ND,NE, HKNOT
57 FORMAT(//1 CURVS test K = ,I2, , ND = ,I5, , NE = ,I5,

```

```

56 CRD1
57 CRD1
58 CRD1
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100 CRD1
101 CRD1
102 CRD1
103 CRD1
104 CRD1
105 CRD1
106 CRD1

1 ' , HKNOT = ,F8.4// DATA... '/'
2 CALL CURV1( K,CB,HKNOT, JG,G, NY,LA,A, LOW,LOWF,JS, T )
3
4 JCON may be either 0 or 1
5 = 0 Factorization is computed after data are read
6 = 1 Factorization is computed as the data are read
7
8 JCON = 0
9 DO 10 I=1,ND
10 READ(5,10030) Y
11 WRITE(6,10030) Y
12 FORMAT(2F10.6)
13 CALL CURV2( JCON, K,CB,HKNOT, LOW,LOWF,T, Y,YO,
14 JG,G,G, NY,LA,A,A, LFLAG )
15 IF(LFLAG.NE.0) GO TO 1001
16 JD = I
17
18 C Compute fit
19 20 CALL CURV3( JCON, K,CB, LA, LOW,LOWF,JS, JG,G,G,
20 NY,LA,A, A,A, LCK,LCM,CPP, MK, LFLAG )
21 IF(LFLAG.NE.0) GO TO 1001
22
23 C Evaluate fit
24
25 WRITE(6,10040)
26 FORMAT(//) Fit evaluation .. '/'
27 DO 39 I=1,NE
28 TX = (I-1)*T/(NE-1)
29
30 C Compute fit
31 E1 = EVAL( K,MK,0.,HKNOT,LCK,CPP(1,1,1), TX )
32 E2 = EVAL( K,MK,0.,HKNOT,LCK,CPP(1,1,2), TX )
33 WRITE(6,10030) E1,E2
34
35 C Print out K, MK, HKNOT, T, and basis coefficients
36 NB = MK+K-2
37 WRITE(6,10050) K,MK,HKNOT, T, NB, (A(I,1),A(I,2),I=1,NB)
38 FORMAT(//1 Spline fit K = ,I2, , MK = ,I3, , HKNOT = ,F10.6CRD1
39 /' Total arc length = ,F10.6//
40 1 X,I3,' Basis coefficients.../(2F10.6)
41 2
42 C Print out piecewise polynomial
43 MKM1 = MK-1
44 DO 40 I=1,NY
45 WRITE(6,10060)
46 FORMAT(// Piecewise polynomial.... ' )
47 DO 40 I=1,MKM1
48 WRITE(6,10070) CPP(J,I,1), J=1,K)
49 FORMAT(8F10.6)
50 STOP
51
52 C ERROR
53 1001 WRITE(6,10080) LFLAG, JD, LOW, JS, T
54 10080 FORMAT( ERROR ,I10// JD,LOW,JS ,3I5, , T = ,F8.2)
55 STOP
56 END

```

```

HRNOT = (BB - AA)/(MK-1)
CALL SETUP( JPER, K, CB, HRNOT, NA, JG, G, NY, LA, A )
C Form normal equations
C
C XL = AA
C LW = 0
C MW = 1.
C DO 40 I=1,ND
C IF( WD(I).GE.0 ) MW = WD(I)
C DX = XD(I) - XL
C IF( (DX.GT.0).AND.(DX.LT.HKNOT) ) GO TO 30
C Compute new LW and HKNOT
C LW = MAX0( 0, MIN0( MK-2, INT( (XD(L)-AA)/HKNOT ) ) )
C XL = AA + LW*HKNOT
C DX = XD(L)-XL
C
C Add into matrices
C 30 CALL ADDON( K, CB, DX, MW, NY, LX, YD( L, 1 ), LW, JG, G, LA, A )
C 40 CONTINUE
C Fold for periodicities
C IF( JPER.NE.0 ) CALL FOLD( K, MK, JG, G, NY, LA, A )
C
C Solve linear systems
C
C CALL FACTOR( 1, N, JG, G, G, NY, LA, A, A, LFLAG )
C Error return if system not positive definite
C IF( LFLAG.NE.0 ) RETURN
C CALL SOLVE( 1, N, JG, G, NY, LA, A, A )
C Periodic extension of coefficients
C IF( JPER.EQ.0 ) GO TO 60
C IF( RW1.LE.0 ) GO TO 60
C DO 50 I=1, NY
C DO 50 I=1, KWI
C A(I+MK-1, I) = A(I, I)
C
C Convert from B-spline to Piecewise polynomial
C
C 60 CALL CONVRT( K, CB, 1, MK, NY, LA, A, LCK, LCM, CPP )
C RETURN
C END
C****END***SPLSQ*****
SPLS 111
SPLS 112
SPLS 113
SPLS 114
SPLS 115
SPLS 116
SPLS 117
SPLS 118
SPLS 119
SPLS 120
SPLS 121
SPLS 122
SPLS 123
SPLS 124
SPLS 125
SPLS 126
SPLS 127
SPLS 128
SPLS 129
SPLS 130
SPLS 131
SPLS 132
SPLS 133
SPLS 134
SPLS 135
SPLS 136
SPLS 137
SPLS 138
SPLS 139
SPLS 140
SPLS 141
SPLS 142
SPLS 143
SPLS 144
SPLS 145
SPLS 146
SPLS 147
SPLS 148
SPLS 149
SPLS 150
SPLS 151
*****
CURVS driver
Incremental Least Squares Splines
Mode 1 Test
29 Aug 74
*****
Externals: from module CURVS: CURV1,CURV2,CURV3
           from module PSQAES: EVAL
Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
"A Real-Time Algorithm for Least Squares Splines and
Its Application in Computer Aided Geometric Design",
research report #29, Department of Computer Science,
Yale University
This program is a driver for the CURVS incremental curve
fitting package. It reads data from a file and prints the
basis coefficients and piecewise polynomial for the fit.
For LFLAG error messages, see module CURVS
C A sample data file:
C...04 10 50 6.5 K,ND,NE,HKNOT (12,215,F10.3)
C... 1. 9. X,Y (2F10.6)
C... 2. 8.5 K is the degree+1 (K le 6)
C... 3. 8. ND is the number of data points
C... 4. 6.5 NE is the number of evaluation points
C... 4.5 5. HKNOT is the knot spacing
C... 5. 3.
C... 4. 2.5
C... 3. 3.
C... 2. 4.
C... 1. 5.
C Resulting Piecewise polynomials (to check operation of program)
C 1.071615 0.733437 0.051666 -0.011503 For X coordinate
C 4.862930 -0.052865 -0.172635 0.012959
C 8.858785 0.167781 -0.279348 0.021057 For Y coordinate
C 3.929716 -0.794753 0.131266 0.004390
DIMENSION Y(2), YO(2)
DATA NY/2/
DIMENSION A(80,2), JG(80), C(6,80)
DIMENSION CPP(6,80,2), CB(6,6)
DATA LCK/6/, LA/80/, LCM/80/
C Fit with Spline
C
C Read K,ND,HKNOT and form matrices
C READ(5,10010) K, ND,NE, HKNOT
C 10010 FORMAT(12,215,F10.3)
C WRITE(6,10020) K, ND,NE, HKNOT
C 10020 FORMAT(/1 CURVS test K = ,12,, ND = ,15,, NE = ,15,

```



```

1      , HKNOT = ,F8.4//
C Initialization
CALL CURV1( K,CB,HKNOT, JG,G, NY,LA,B, LOW,LOWF,JS, T )
C
DO 30 I=1,ND
  JD = I
  READ(5,10030) Y(1),Y(2)
  FORMAT(2F10.6)
C Add contribution to matrices
CALL CURV2( -1, K,CB,HKNOT, LOW,LOWF,T, Y,XO, JG,G,G,
  NY,LA,B,B, LFLAG )
  IF(LFLAG.NE.0) GO TO 1001
  IF(1.LT.K) GO TO 30
C Compute fit
CALL CURV3( -1, K,CB, LA, LOW,LOWF,JS, JG,G,G,
  NY,LA,B,B, LCK,LCK,CPP, MK, LFLAG )
  IF(LFLAG.NE.0) GO TO 1001
C Print results
MKM1 = MK-1
DO 20 L=1,NY
  WRITE(6,10035) L, (A(J,L), J=1,JS)
  FORMAT(/, BASIS COEFFICIENTS..., I2/(6F12.5))
10035
  WRITE(6,10040) L
  FORMAT(/, Piecewise polynomial..., I2)
10040
  DO 20 M=1,MKM1
    WRITE(6,10050) (CPP(J,M,L), J=1,K)
    FORMAT(6F12.5)
  20
  CONTINUE
C ERROR
1001 WRITE(6,10060) LFLAG, JD, LOW, JS, T
10060 FORMAT(' ERROR ',I10/' JD,LOW,JS ',3I5,' T = ',F8.2)
STOP
END

```

```

CRD2 56
CRD2 57
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CRD2 59
CRD2 60
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CRD2 63
CRD2 64
CRD2 65
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CRD2 67
CRD2 68
CRD2 69
CRD2 70
CRD2 71
CRD2 72
CRD2 73
CRD2 74
CRD2 75
CRD2 76
CRD2 77
CRD2 78
CRD2 79
CRD2 80
CRD2 81
CRD2 82
CRD2 83
CRD2 84
CRD2 85
CRD2 86
CRD2 87
CRD2 88
CRD2 89

*****
* CURVS driver
* Incremental Least Squares Splines
* Mode 3 Test
* 29 Aug 74
*****
Externals: from module CURV1,CURV2,CURV3,CURV4
           from module PSQARS: EVAL

Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
"A Real-time Algorithm for Least Squares Splines and
Its Application in Computer Aided Geometric Design",
research report #29, Department of Computer Science,
Yale University

This program is a driver for the CURVS incremental curve
fitting package. It reads data from a file and prints the
basis coefficients and piecewise polynomial for the fit.

For LFLAG error messages, see module CURVS

C A sample data file:
C... 4 11 25 1.2567 K,ND,NE,HKNOT (12,215,F10.3)
C... 1.000000 0. X,Y
C... 0.809017 0.587785 K is the degree-1 (K LE 6)
C... 0.309017 0.951057 ND is the number of data points
C... -0.809017 0.587785 NE is the number of evaluation points
C... -1.000000 0. HKNOT is the knot spacing
C... -0.809017 -0.587785
C... -0.309017 -0.951057
C... 0.309017 -0.951057
C... 0.809017 -0.587785
C... 1.000000 0.

C Resulting basis coefficients (to check operation of program)
BASIS COEFFICIENTS... 1
0.28908 1.33392 0.37554 -1.09604 -1.01319 0.50419
1.33095 0.11767
BASIS COEFFICIENTS... 2
-1.29037 0.00864 1.25699 0.73147 -0.84184 -1.21183
0.13368 1.29697

DIMENSION Y(2), YO(2)
DATA NY/2/
DIMENSION A(80,2), ALIB(80,2), B(80,2)
DIMENSION JG(80), G(11,80), GLD(6,80)
DIMENSION CPP(6,80,2), CB(6,6)
DATA LCK/6/, LA/80/, LCK/80/

C Fit with Spline
C Read K,ND,HKNOT and form matrices

```

```

CRD3 1
CRD3 2
CRD3 3
CRD3 4
CRD3 5
CRD3 6
CRD3 7
CRD3 8
CRD3 9
CRD3 10
CRD3 11
CRD3 12
CRD3 13
CRD3 14
CRD3 15
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CRD3 54
CRD3 55

```

```

10010 READ(5,10010) K, ND, NE, HKNOT
      FORMAT(12,2I5,F10.3)
10020 WRITE(6,10020) K, ND, NE, HKNOT
      FORMAT(/,'1 CURVS test K = ',12,' ND = ',15,' NE = ',15,
1      ' HKNOT = ',F8.4/)
C Initialization
C CALL CURV1( K,CB,HKNOT, JG,G, NY,LA,B, LOW,LOWF,JS, T )
      DO 10 I=1,ND
          JD = I
          READ(5,10030) Y(1),Y(2)
          FORMAT(2F10.6)
C Add to matrices
1          CALL CURV2( 1, K,CB,HKNOT, LOW,LOWF,T, Y,YO, JG,G,GLD,
2          NY,LA,B,ALIB,A, LCK,LCM,CPP, MK, LFLAG )
C Compute fit (open curve)
3          IF(LFLAG.NE.0) GO TO 1001
          IF(1.LT.K) GO TO 10
C Compute final fit (closed curve)
4          CALL CURV3( 1, K,CB, LA, LOW,LOWF,JS, JG,G,GLD,
5          NY,LA,B,ALIB,A, LCK,LCM,CPP, MK, LFLAG )
          IF(LFLAG.NE.0) GO TO 1001
C Print out fit
6          DO 5 I=1,NY
7              WRITE(6,10035) I, (A(J,L), J=1,JS)
8              FORMAT(/,'BASIS COEFFICIENTS...',I2/(6F12.5))
9          CONTINUE
C Compute final fit (closed curve)
10         CALL CURV4( K,CB, LOW,JS, JG,G, NY,LA,A,B,
11         LCK,LCM,CPP, MK, LFLAG )
          IF(LFLAG.NE.0) GO TO 1001
C Evaluate fit
12         WRITE(6,10040)
13         FORMAT(/,' Fit evaluation .. '/')
14         C = (LOW+1)*HKNOT
          DO 20 I=1,NE
              TX = (I-1)*C/(NE-1)
              E1 = EVAL( K,MK,0.,HKNOT,ICK,CPP(1,1,1), TX )
              E2 = EVAL( K,MK,0.,HKNOT,ICK,CPP(1,1,2), TX )
              WRITE(6,10045) E1,E2
              FORMAT(2F12.5)
          STOP
C ERROR
1001 WRITE(6,10050) LFLAG, JD, LOW, JS, T
10050 STOP
      STOP
      END

```

```

CRD3 56
CRD3 57
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CRD3 102
CRD3 103

```

```

*****
* * Module CURVS
* * Fast Least Squares Spline Curve Fitting
* * 28 Aug 74
*****
C Entries: CURV1, CURV2, CURV3, CURV4
C External: from module PSQAKS: SETUP,ADDON,EXPAND,FOLD,CONVRT
           from module ENSOLV: FACTOR,SOLVE
C Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
           "A Real-Time Algorithm for Least Squares Splines and
           Its Application in Computer Aided Geometric Design",
           research report #29, Department of Computer Science,
           Yale University
C The CURVS module is designed for incremental fitting of curves
C in two or more dimensions. The subroutines minimize storage and
C execution time while enabling evaluation of the fit at any time
C during acquisition of data. There are four routines: CURV1 for
C initialization, CURV2 for data point processing, CURV3 for open
C curve fit computation, and CURV4 for closed curve fit computation.
C Both CURV3 and CURV4 return the resulting spline as a piecewise
C polynomial which may be evaluated efficiently using EVAL in the
C PSQAKS module.
C This module may be used in three basic modes:
C 1) Sequential: The fit is not required until all data are acquired.
           Call CURV1 for initialization; call CURV2 once for each data
           point; and call CURV3 or CURV4 to compute the fit.
           G and GLD may use the same storage.
           A, B, and ALIB may use the same storage.
           The control variable JCON may be set to 0 or 1. If JCON=1,
           the factorization GLD is computed in CURV2 while the data
           are being acquired; and if JCON=0, the entire factorization
           is computed in CURV3.
C 2) Incremental: The fit will be computed several times before the
           complete set of data is acquired.
           Call CURV1 for initialization and CURV2 once for each data point.
           CURV3 may be called at any time to compute the fit.
           G and GLD may use the same storage.
           A, B and ALIB may use the same storage.
           If G and GLD use the same storage, set JCON=1,
           otherwise set JCON=-1.
           For an exact fit, (but requiring more computation time as the
           curve grows), set JLAG = LA. For an approximate, but nearly
           exact fit, set JLAG = 2*K (approximately).
C 3) Incremental closed: The fit will be computed several times
           before the complete set of data is acquired; and after the
           complete set of data is acquired, a closed curve fit
           will be computed. Call CURV1 for initialization and CURV2
           once for each data point. CURV3 may be called at any time
           to compute the open curve fit. To compute the closed

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CRVS 218
CRVS 219
CRVS 220

RETURN
END
C***CURV2*****
SUBROUTINE CURV2(JCON, K,CB,HKNOT, LOW,LOWF,T, Y,YO, JG,G,GLD,
1 NDIM,LA,B,ALIB, LFLAG )
DIMENSION CB(K,K), JG(LA),G(1),GLD(1)
DIMENSION B(LA,NDIM),ALIB(LA,NDIM)
DIMENSION Y(NDIM), YO(NDIM)
LFLAG = 0
C Arc length
O = 0.
DO 10 I=1,NDIM
IF(T.LT.0.) GO TO 10
O = O + (Y(I)-YO(I))**2
10 YO(I) = Y(I)
T = T + SQRT(O)
IF(T.LT.0) T = 0.
C Compute interval
JO = 1
IF(JCON.LT.0) JO = K+1
DX = T-LOW*HKNOT
IF(DX.LT.HKNOT) GO TO 30
LOW = LOW+1
IF( (LOW+K-JO-1).LE.LA ) GO TO 20
LFLAG = 1
RETURN
C Add to matrix
30 JOG = JG(JO)+JO
CALL ADDON(K,CB,DX,1,NDIM,1,Y, LOW, JG,G(JOG),LA,B(JO,1))
C Factor if necessary and requested
IF(JCON.EQ.0) RETURN
IF(LOWF.GE.LOW) RETURN
CALL FACTOR( LOWF+1,LOW, JG,G(JOG),GLD, NDIM,LA,B(JO,1),ALIB,
1 LFLAG )
LOWF = LOW
RETURN
END
C***CURV3*****
SUBROUTINE CURV3( JCON, K,CB, JLAG, LOW,LOWF,JS, JG,G,GLD,
1 NDIM,LA,B,ALIB,A, LCK,LCM,CPP, MK, LFLAG )
DIMENSION CB(K,K), JG(LA),G(1),GLD(1)
DIMENSION B(LA,NDIM),ALIB(LA,NDIM), A(LA,NDIM)
DIMENSION CPP(LCK,LCM,NDIM)
JS = LOW+K
C Factor remainder of matrix
JO = 1
IF(JCON.LT.0) JO = K+1
JOG = JG(JO)+JO
LOWF = LOW
IF(JCON.EQ.0) LOWF = 0
IF((LOWF+1).LT.JS) CALL FACTOR(LOWF+1,JS-1, JG,G(JOG),GLD,
1 NDM,LA,B(JO,1),ALIB, LFLAG )
IF(LFLAG.NE.0) RETURN
CALL FACTOR(JS,JS,JG,G(JOG),GLD, NDIM,LA,B(JO,1),ALIB, NFLAG)
IF(NFLAG.NE.0) JS = JS - 1

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CRVS 221
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CRVS 256

MK = JS - K + 2
C Solve
JLOW = 1 + MAX0( 0, MIN0( JS, LOWF-JLAG ) )
CALL SOLVE( JLOW,JS, JG,GLD, NDM,LA,A,ALIB )
C Convert
CALL CONVRT( K,CB, JLOW,MK, NDM,LA,A, LCK,LCM,CPP )
RETURN
END
C***CURV4*****
SUBROUTINE CURV4( K,CB, LOW,JS, JG,G, NDM,LA,A,B,
1 LCK,LCM,CPP, MK, LFLAG )
DIMENSION CB(1), JG(LA), G(1), A(LA,NDIM), B(LA,NDIM)
DIMENSION CPP(LCK,LCM,NDIM)
MK = LOW+2
JS = LOW+K
LFLAG = 2
IF(MK.LE.K) RETURN
C Expand and fold matrix G
CALL EXPAND( K, MK, JG,G )
CALL FOLD( K, MK, JG,G, NDM,LA,B )
C Factor G
JS = LOW+1
CALL FACTOR( 1,JS, JG,G,G, NDM,LA,B,A, LFLAG )
IF(LFLAG.NE.0) RETURN
C Solve
CALL SOLVE( 1,JS, JG,G, NDM,LA,A,A )
IF(K.LE.1) GO TO 20
C Extend basis coefficients
DO 10 I=1,NDIM
DO 10 I=2,K
A(I+MK-2,I) = A(I-1,I)
10
C Convert
20 CALL CONVRT( K,CB, 1,MK, NDM,LA,A, LCK,LCM,CPP )
RETURN
END
C***END***CURVS*****

```

```

C curve fit, call CURV4 after acquisition of the complete
C set of data. G, GLD, A, ALIB, and B must use different
C storage. Observe that the array G used for a closed curve is
C DIMENSIONED nearly twice as large as that for an open curve.
C JCON should be set to -1 and JLAG is chosen as in (2).
C
C To achieve maximum efficiency, the user may wish to modify
C the organization of the modules CURVS, PSQARS, and IPSOLV.
C The current organization is designed for clarity rather than
C for efficiency. For example, the subroutine call to CURV2
C may require more time than the computation in CURV2 itself.
C By incorporating some of the arguments in common blocks and
C eliminating some subroutine calls entirely, the user may speed
C up these routines considerably.
C The same considerations apply to evaluation of the fit.
C The user interested in efficiency should not call the function
C EVAL to evaluate a spline; instead that calculation should be
C performed with in-line code.
C
C*****Subroutines and Calling Sequences*****
C
C 1) Unless otherwise indicated, variable types are
C IMPLICIT REAL(A-H, O-Z)
C IMPLICIT INTEGER(I-N)
C
C 2) Value variables (V) pass a value to the subroutine and must
C have been set by the user or by a previously called subroutine.
C Result variables (R) return results to the calling subprogram.
C
C**CURV1: Setup subroutine for CURVS module
C CALL CURV1( K,CB,HKNOT, JG,G,NDIM,LA,B, LOW,LOWF,JS,T )
C
C V K
C R CB(K**2) - Order of spline
C G(LA,NDIM) - The piecewise polynomial representation
C for the order k B-spline basis on uniform
C knots with spacing HKNOT
C V HKNOT - Knot spacing: if T is the total arc length,
C then LA must be GE T/HKNOT
C V LA - Dimension of JG,G,B
C R JG(LA) - Index array for gramman
C R G(K,LA) - Gramman (the profile of the lower triangle is stored)
C G(2*K-1,LA) The first set of dimensions is sufficient for
C routines CURV1,CURV2,CURV3; but if CURV4 is to be
C called, then G must have the second dimensions
C NOTE: the indicated dimensions are for the
C purpose of allocating storage only
C the array is actually addressed as follows:
C G(I,J) = G( JG(I) + J, I ) for I le J
C
C V NDIM - Dimension of curve
C R B(LA,NDIM) - Right hand sides (NDIM of them, LA long)
C R LOW - Interval variable, (LOW*HKNOT) LE T LT ( (LOW+1)*HKNOT)
C R LOWF - Value of LOW at last factorization of A
C R JS - Index of highest valid basis coefficient
C R T - LOW,LOWF,JS are initialized to 0
C R T - Accumulated arc length (initialized to -1.)

```

```

C**CURV2
C 1) Computes the arc length T to data point Y(NDIM)
C 2) Adds contribution of data point Y(NDIM) to G and B
C 3) Factors G and performs forward solve on B, both up to row LOW
C
C CALL CURV2( JCON, K,CB,HKNOT, LOW,LOWF,T, Y,YO,
C JG,G,GLD, NDIM,LA,B,ALIB, LFLAG )
C
C V JCON - Control flag (set mode of operation)
C 0 - Do not factor matrix G, (mode 1)
C -1 - Factor G; and CURV3 will be called (mode 2)
C 1 - Factor G (mode 3)
C V K,CB,HKNOT - As previously defined
C V JG,NDIM,LA
C V G,B
C VR LOW,LOWF,T
C
C V Y(NDIM) - Coordinates of point on curve
C VR YO(NDIM) - Previous coordinates
C R GLD(K,LA) - Factorization of G (if requested) to row LOW
C R ALIB(LA) - Forward solve of B to element LOW
C R LFLAG - Error code (0, OK; 1, LA too small; -1, G singular)
C
C**CURV3: Computes fit
C CALL CURV3( JCON, K,CB, JLAG, LOW,LOWF,JS, JG,G,GLD,
C NDIM,LA,B,ALIB,A, LCK,LCM,CPP, MK, LFLAG )
C
C V JCON,K,CB,HKNOT - As previously defined
C V JG,NDIM,LA
C V LOW,G,B
C VR GLD,ALIB,LOWF,JS,
C R LFLAG
C
C V JLAG - Number of intervals lower than LOWF to be backsolved
C (for exact solution JLAG = LA)
C R A(LA,K) - Basis coefficients for resulting spline
C R CPP(LCK,LCM,NDIM) - Piecewise polynomial for resulting spline
C R MK - Number of knots used
C
C**CURV4: Computes closed curve fit
C CALL CURV4( K,CB, LOW, JG,G, NDIM,LA,B,A, LCK,LCM,CPP,MK, LFLAG)
C
C V K,CB - As previously defined
C V NDIM,LA,LOW
C VR JG,G,B
C R A,CPP,MK,LFLAG
C
C**CURV1*****
C SUBROUTINE CURV1( K,CB,HKNOT, JG,G,NDIM,LA,B, LOW,LOWF,JS, T )
C DIMENSION CB(K,K), JG(LA), G(L1), B(LA,NDIM)
C CALL SETUP3( 0,K,CB,HKNOT, LA, JG,G, NDIM,LA,B )
C LOW = 0
C LOWF = 0
C JS = 0
C T = -1.

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C *****
C * Module PSQMS
C * Least Squares Polynomial Splines
C * Primitives
C * 22 Sep 74
C *****
C C Entries: SETUP, ADDON, EXPAND, FOLD, CONVRT, EVAL
C External: none beyond FORTRAN IV library
C These subroutines are to be used in the incremental formation
C of the spline normal equations
C
C Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
C "A Real-Time Algorithm for Least Squares Splines and
C Its Application in Computer Aided Geometric Design",
C research report #29, Department of Computer Science,
C Yale University
C *****Subroutines and Calling Sequences*****
C
C 1) Unless otherwise indicated, variable types are
C IMPLICIT REAL(A-H, O-Z)
C 2) Value variables (V) pass a value to the subroutine
C Result variables (R) return results to the calling subprogram
C
C**SETUPS
C This subroutine must be called before using ADDON or CONVRT
C 1) Sets up basis function array CB
C 2) Sets up index array JG
C 3) Zeros parts of arrays G and B
C
C CALL SETUP( JPER, K,CB,HMESH, NSET, JG,G, NY,LB,B )
C
C V JPER - Periodic flag, if non-zero setup for periodic splines
C V K - Order of spline
C R CB(K**2) - The piecewise polynomial representation
C for the order k B-spline basis on uniform
C knots with spacing HMESH
C
C V HMESH - Knot spacing
C V NSET - Number of entries to be set in JG, G, and B
C V LB,NY - Dimensions of B
C R G(*) - Gram matrix (profile only is stored)
C G(I,J) = G( JG(I) + J ) for I le J
C R JG(LB) - Index array for Gram matrix
C R E(LB,NY) - Right hand sides (NY of them, LB long)
C
C RESTRICTIONS: K must be LE 6, NSET must be LE LB,
C and dim(G) must be GE (NSET*K - K*(K-1)/2)
C
C**ADDON
C Adds contribution of measurements at points ( DX*LOW*HMESH, Y )
C to the matrices G and B
C
C CALL ADDON( K,CB, DX,W, NY,LY,Y, LOW, JG,G, LB,B )

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C V LOW, DX - Data point is at DX + LOW*HMESH + XLEFT
C where XLEFT is the left hand end point of the interval
C V W - Data weight
C V NY,LY - Dimensions of Y
C V Y(LY,NY) - Data values (NY of them spaced at intervals LY in Y)
C
C V JG(LB),LB,CB(K**2) - As previously defined
C VR G(*),B(LB)
C RESTRICTIONS: SETUPS must have been called first with NSET GE (LOW*K)
C and LB must be GE (LOW*K)
C
C**EXPAND
C Expands matrix G stored as a K wide band into the shape needed
C for the periodic matrix
C
C CALL FOLD( K, MK, JG,G )
C
C V K,JG(MK+K-2) - As previously defined
C VR G(*)
C RESTRICTIONS: SETUPS must have been called first with NSET GE (MK+K-2)
C and dim(G) must be GE (2*K-1)*(MK+K-2)
C
C**FOLD
C Folds non-periodic matrices to create periodic matrices. EXPAND
C or the equivalent routine must have been called first
C
C CALL FOLD( K, MK, JG,G, NY,LB,B )
C
C V K,JG(LB),NY,LB - As previously defined
C VR G(*),B(LB)
C RESTRICTIONS: (MK+K-2) must be LE LB,
C and dim(G) must be LE (MK+K-2)*(2*K-1)
C
C**CONVRT
C Converts basis coefficients A into piecewise polynomial
C coefficients CPP
C
C CALL CONVRT( K,CB, JLOW,JHIGH, NY,LA,A, LCK,LCM,CPP )
C
C V JLOW,JHIGH - Intervals for which pp is to be computed
C V LA,NY - Dimensions of A
C V A(LA,NY) - B-spline basis coefficients
C V LCK,LCM,NY - Dimensions of CPP
C R CPP(LCK,LCM,NY) - Array of piecewise polynomial coefficients
C
C K,CB - As previously defined
C RESTRICTIONS: SETUPS must have been called first, K must be LE LCK,
C (JHIGH-1) must be LE LCM
C**EVAL
C Evaluates spline at XX from pp representation

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C C VALUE = EVAL( K,MK,XLEFT,HMESH, LCK, CPP, XX )
C C
C V XX
C V K,MK,XLEFT,HMESH,LCK, CPP - As previously defined
C C
C**SETUPS*****
SUBROUTINE SETUPS( JPER, K,CB,HMESH, NSET, JG,G, NY, LB, B )
DIMENSION CB(K,K), JG(LB), G(1), B(LB,NY)

C CB is an array containing columns of the
C K coefficients for each of the K different piecewise polynomials
C of the K order B-spline basis functions on an integer mesh.
C The entries for order K begin at CBM( JCBM(K)+1 ).
C K must be less than KWAX = 6.

DIMENSION JCBM(6), CBM(91), CBM1(30), CBM2(25), CBM3(36)
EQUIVALENCE (CBM,CBM1), (CBM(31),CBM2), (CBM(56),CBM3)
DATA JCBM/ 0,1,5,14, 30,55/
DATA CBM1/
1 1., 1., -1., 0., +1.,
2 5.0000000E-01, -1.0000000E+00, 5.0000000E-01,
3 5.0000000E-01, 1.0000000E+00, -1.0000000E+00,
4 0.0000000E-01, 0.0000000E-01, 5.0000000E-01,
5 1.6666667E-01, -5.0000000E-01, 5.0000000E-01, -1.6666667E-01,
6 6.6666666E-01, 0.0000000E-01, -1.0000000E+00, 5.0000000E-01,
7 1.6666667E-01, 5.0000000E-01, 5.0000000E-01, -5.0000000E-01,
8 0.0000000E-01, 0.0000000E-01, 0.0000000E-01, 1.6666667E-01/
DATA CBM2/
1 4.1666667E-02, -1.6666667E-01, 2.5000000E-01, -1.6666667E-01,
2 4.1666667E-02, 4.5833333E-01, -5.0000000E-01, -2.5000000E-01,
3 5.0000000E-01, -1.6666667E-01, 4.5833333E-01, 5.0000000E-01,
4 -2.5000000E-01, -5.0000000E-01, 2.5000000E-01, 4.1666667E-02,
5 1.6666667E-01, 2.5000000E-01, 1.6666667E-01, -1.6666667E-01,
6 0.0000000E-01, 0.0000000E-01, 0.0000000E-01, 0.0000000E-01,
7 4.1666667E-02/
DATA CBM3/
1 8.3333333E-03, -4.1666667E-02, 8.3333333E-02, -8.3333333E-02,
2 4.1666667E-02, -8.3333333E-03, 2.1666667E-01, -4.1666666E-01,
3 1.6666667E-01, 1.6666667E-01, -1.6666667E-01, 4.1666667E-02,
4 5.0000000E-01, 0.0000000E-01, -5.0000000E-01, 0.0000000E-01,
5 2.5000000E-01, -8.3333333E-02, 2.1666667E-01, 4.1666666E-01,
6 1.6666667E-01, -1.6666667E-01, -1.6666667E-01, 8.3333333E-02,
7 8.3333333E-03, 4.1666667E-02, 8.3333333E-02, 8.3333333E-02,
8 4.1666667E-02, -4.1666667E-02, 0.0000000E-01, 0.0000000E-01,
9 0.0000000E-01, 0.0000000E-01, 0.0000000E-01, 8.3333333E-03/

C Set up index array JG
C C
KMI = K-1
JG(1) = 0
DO 10 I=2,NSET
JG(I) = JG(I-1) + MIN0(I,K) - 1
IF (JPER.EQ.0) GO TO 40
IF (KMI.LE.0) GO TO 40

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PSOR 165

M = 0
JF = NSET-2*K+2
DO 20 I=1,KMI
M = M + MAX0( 0, I+JF-K )
JG(I+JF) = JG(I+JF) + M
JF = JF+KMI
DO 30 I=1,KMI
JG(I+JF) = JG(I+JF) + M
Zero array G, and NY columns of B
DO 40 I=1,NY
DO 50 I=1,NSET
B(I,L) = 0.
NN = JG(NSET) + NSET
DO 60 I=1,NN
G(I) = 0.

C Set up TABLES common
C C
F = 1.
DO 80 I=1,K
IJ = JCBM(K) + I
DO 70 J=1,K
CB(I,J) = CBM(IJ)*F
IJ = IJ+K
F = F/HMESH
RETURN
END
C**ADDON*****
SUBROUTINE ADDON( K,CB, DX,W, NY,LY,Y, LOW, JG,G, LB,B )
DIMENSION CB(K,K), Y(LY,NY), JG(LB), G(1), B(LB,NY)
DIMENSION VB(L0)
DO 20 I=1,K
VV = 0.
DO 10 J=1,K
VV = DX*VV + CB(K-J+1,I)
VB(I) = VV
C Add into matrix G
VBW = VB(I)*W
II = JG(I+LOW)+LOW
DO 40 J=1,I
G(II+J) = G(II+J) + VBW*VB(J)
C Add into B
DO 60 L=1,NY
B(I+LOW,L) = B(I+LOW,L) + Y(I,L)*VBW
CONTINUE
RETURN
END
C**EXPAND*****
SUBROUTINE EXPAND( K, MK, JG,G )
DIMENSION JG(1), G(1)
KMI = K-1
IF (KMI.LE.0) RETURN

```

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PSQR 296

DO 30 I=1,NY
DO 30 J=JLOW,NM
DO 20 J=1,K
DD = 0.
DO 10 I=1,K
DD = DD + A(I+M-1,L)*CB(J,I)
CPP(J,M,L) = DD
CONTINUE
RETURN
END
C***EVAL*****PSOARS*****
FUNCTION EVAL( K,MK,XLEFT,HMESH, LCK, CPP, XX )
DIMENSION CPP(LCK,1)
LOW = MAX0( 0, MIN0( MK-2, INT( (XX-XLEFT)/HMESH ) ) )
DX = XX - (XLEFT + LOW*HMESH)
EVAL = 0.
DO 10 I=1,K
EVAL = EVAL*DX + CPP(K-I+1,LOW+1)
RETURN
END
C***END***PSOARS*****

```

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PSQR 275

C
Expand and Copy
N = 0
LM = K*K*M1
DO 10 I=1,KM1
LM = LM + MIN0( I+JF2, K )
N = N + MAX0( 0, I+JF2-K )
JG(I+JF2) = JG(I+JF2) + M
IF(M.LE.0) RETURN
DO 20 I=1,KM1
JG(I+JF1) = JG(I+JF1) + M
JGN = JCO+M
DO 30 I=1,LM
G(JCN-I+1) = G(JCO-I+1)
II = JCN-LM
DO 70 I=1,KM1
JGT = JG(I+JF2)
LL = JF2-I-K
IF(LL.LE.0) GO TO 50
DO 40 J=1,LL
G(JGT+J) = 0.
LN = MIN0( K, JF2+I )
JGU = JGT+JF2+I-LN
DO 60 J=1,LN
G(JGU+J) = G(II+J)
II = II+LN
RETURN
END
C***FOLD*****
SUBROUTINE FOLD( K, MK, JG,G, NY, LB, B )
DIMENSION JG(LB), G(1), B(LB,NY)
K1 = K-1
IF(K1.LE.0) RETURN
JF1 = MK-1
JF2 = MK-K
C Fold
DO 30 I=1,KM1
JGD1 = JG(I)
JGS1 = JG(I+JF1)+JF1
JGD2 = JG(JF2+I)
DO 10 L=1,NY
B(I,L) = B(I,L) + B(I+JF1,L)
DO 20 J=1,I
JGS2 = JG(JF1+J) + JF2+I
G(JGD2+J) = G(JGD2+J) + G(JGS2)
G(JGD1+J) = G(JGD1+J) + G(JGS1+J)
CONTINUE
RETURN
END
C***CONVRT*****
SUBROUTINE CONVERT( K,CB, JLOW,JHIGH, NY,LA,A, LCK,LCM, CPP )
DIMENSION CB(K,K), A(LA,NY), CPP(LCK,LCM,NY)
NM = JHIGH-1

```

```

C *****
C * Incremental Symmetric Envelope Linear Equation Solver
C * 29 Aug 74
C *****
C Entries: FACTOR, SOLVE
C Externals: none beyond FORTRAN IV library
C
C Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
C "A Real-Time Algorithm for Least Squares Splines and
C Its Application in Computer Aided Geometric Design",
C Research report #29, Department of Computer Science,
C Yale University
C
C The module ENSOLV computes solutions to linear systems  $A^*X = B$ 
C where A is an  $N \times N$  positive definite matrix of which half of the
C envelope is stored; B is an arbitrary  $N \times NB$  matrix, and X is the
C  $N \times NB$  solution matrix.
C The system is solved in three steps:
C a) Factoring the matrix A as  $L^*D^*(L$  transpose), here called ALD,
C where L is lower triangular with unit diagonal and D is diagonal
C b) Solving the triangular system  $L^*ALIB = B$  for ALIB
C c) Solving the triangular system  $D^*(L$  transpose) $*X = ALIB$  for X
C
C The matrix A and its factorization ALD are stored in one
C dimensional arrays and accessed through the index array JA:
C  $A(I,J) = A(JA(I) + J)$  for  $I \leq J$ 
C Only the lower triangle of the envelope of A and ALD is stored since
C all other elements are zero. For example, consider the lower
C triangle of matrix A, its envelope, and its index array JA:
C
C A: X envelope(A): E JA: 0
C XY EF JA: 1
C 0XX EF JA: 2
C 00XX EF JA: 2
C X00XX FEFE JA: 6
C 00XX FEFE JA: 9
C
C For a given row of the lower triangle, the envelope is composed
C of those elements in the row of same or higher column index than
C the nonzero element of lowest column index. To compute the index
C vector JA, use the formula:
C
C  $JA(I) = 0$ 
C  $JA(I) = JA(I-1) - 1 +$  (number of elements in the envelope for row I)
C
C To solve a simple linear system,  $A^*X = B$ 
C CALL FACTOR( 1,N, JA,A,A, 1,0,B,B, LFLAG )
C CALL SOLVE ( 1,N, JA,A, 1,0,X,B )
C LFLAG is a status code, zero indicates successful completion.
C (in this example the factorization LD replaces the contents of A
C and (L inverse) $*B$  replaces the contents of B )
C
C ENSOLV may also provide incremental factorizations and solutions.
C Since the factorization of A is computed row by row, no element of

```

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45 ENSL 100
46 ENSL 101
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54 ENSL 109
55 ENSL 110

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C ALD in a given row depends on any element of A in subsequent rows;
C consequently the factorization for the first 0 rows of a matrix
C may be computed without knowledge of the values for rows 0+1, 0+2,.....
C For example:
C
C a) COMPUTE K rows of the factorization LD (which replace the
C corresponding rows of A) and K elements of (L inverse) $*B$ 
C (which replace the corresponding elements of B)
C CALL FACTOR( 1,K, JA,A,A, 1,0,B,B, LFLAG )
C b) COMPUTE K additional rows
C CALL FACTOR( K+1,2*K, JA,A,A, 1,0,B,B, LFLAG )
C c) COMPUTE K additional rows
C CALL FACTOR( 2*K+1,3*K, JA,A,A, 1,0,B,B, LFLAG )
C
C The array A will contain its factorization LD in the first 3*K
C rows and subsequent rows will be unchanged.
C *****Subroutines and Calling Sequences*****
C 1) Unless otherwise indicated, variable types are
C IMPLICIT REAL(A-H, O-Z)
C IMPLICIT INTEGER(I-N)
C 2) Value variables (V) pass a value to the subroutine
C result variables (R) return results to the calling subprogram
C
C**FACTOR
C 1) Factors A = L*D*(L transpose)
C 2) Forward solves ALIB = (L inverse) $*B$ 
C
C CALL: CALL FACTOR( JBOT,JTOP, JA,A,ALD, NB,LB,B,ALIB, LFLAG )
C
C V JBOT, JTOP- Range over which envelope solution is carried out
C V A(*) - The envelope of the matrix A is stored in row order
C V JA(LB) - Index array for A. To get an element:
C  $A(I,J) = A(JA(I) + J)$  for  $I \leq J$ .
C The number of elements of row I belonging to the
C envelope is  $JA(I) - JA(I-1) + 1$ .
C Length of matrix =  $JA(N) + N$  for  $N \times N$  matrix.
C V B(LB,NB) - Right hand sides (NB of them, LB long)
C R ALD(*) - L D LT decomposition of A
C K ALIB(LB,NB)-L inverse times B
C
C R LFLAG is -1 if A is NOT positive definite, 0 otherwise
C A and ALD may use the same storage
C B and ALIB may use the same storage
C**SOLVE
C Back solves  $X = (LT$  inverse) $*(D$  inverse) $*ALIB$ 
C CALL SOLVE( JBOT,JTOP, JA,ALD, NB,LB,X,ALIB )
C V JBOT,JTOP,NB,LB - As previously defined

```



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1001 LFLAG = -1
RETURN
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C ALD(*), ALIB(*)
C R X(LB,NB) - Solution vectors (NB of them, LB long)
C
C X and ALIB may use the same storage
C
C**FACTOR*****
SUBROUTINE FACTOR( JBOT,JTOP, JA,A,ALD, NB,LB,B,ALIB, LFLAG )
DIMENSION JA(1), A(1), ALD(1), B(1), ALIB(1)
EPS is the single precision rounding error (here for the PDP10)
DMA EPS/1.E-7/
LFLAG = 0
C Factorization
DO 110 K=JBOT,JTOP
JAK = JA(K)
IMAX = K-1
IF (IMAX.LE.0) GO TO 50
IMIN = MAX0( 1, K-JA(K)+JA(K-1) )
IF (IMIN.GE.K) GO TO 50
C Off diagonal elements
DO 40 I=IMIN,IMAX
JAK = I-1
AIK = A(JAK+I)
IF (OMAX.LE.0) GO TO 40
JAI = JA(I)
JMIN = MAX0( I-JA(I)+JA(I-1), IMIN )
IF (JMIN.GE.I) GO TO 40
DO 30 J=JMIN,JMAX
AIK = AIK - ALD(JAI+J) * ALD(J+JAK)
30 ALD(I+JAK) = AIK
40 ALD(I+JAK) = AIK
C Diagonal elements
50 ALKK = A(K+JAK)
IF (IMAX.LE.0) GO TO 70
IF (IMIN.GE.K) GO TO 70
DO 60 I=IMIN,IMAX
JAI = JA(I)
AIK = ALD(I+JAK)
ALIK = AIK*ALD(I+JAI)
ALD(I+JAK) = ALIK
60 ALKK = ALKK - AIK*ALIK
C Check sign of diagonal element (including rounding error)
70 IF (ALKK.LE.ANAXI(0.),(K-1)*EPS*A(K+JAK)) GO TO 1001
ALD(K+JAK) = 1./ALKK
C Forward solve
II = 0
DO 100 J=1,NB
BK = B(K+II)
IF (IMAX.LE.0) GO TO 90
IF (IMIN.GE.K) GO TO 90
DO 80 I=IMIN,IMAX
BK = BK - ALD(JAK+I)*ALIB(I+II)
80 ALIB(K+II) = BK
90 II = II+LB
100 CONTINUE
110 RETURN
C Error return: not positive definite ( Di < 0 )

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Appendix II: PP-Representation of the B-Splines

For t satisfying

$$(II.1) \quad i \cdot h \leq t < (i+1) \cdot h$$

and with

$$(II.2) \quad \delta = t - ih$$

the B-splines may be written as

$$(II.3) \quad N_{i+1,k}(t) = \sum_{\ell=0}^{k-1} C_{1\ell}^N \delta^\ell$$

$$\vdots$$
$$N_{i+k,k}(t) = \sum_{\ell=0}^{k-1} C_{k\ell}^N \delta^\ell$$

$$N_{j,k}(t) = 0 \quad \text{for } j \leq i \text{ or } j > i+k.$$

A table of the coefficients C_{ij}^N follows. This table was generated by rationalizing the numbers generated by the program PPBAS. If higher order splines are desired, PPBAS may be used to extend this table.

Order	Coefficients					
	δ^0	δ^1	δ^2	δ^3	δ^4	δ^5
1	1					
2	1	-1				
	0	1				
3	1/2	-1	1/2			
	1/2	1	-1			
	0	0	1/2			
4	1/6	-1/2	1/2	-1/6		
	2/3	0	-1	1/2		
	1/6	1/2	1/2	-1/2		
	0	0	0	1/6		
5	1/24	-1/6	1/4	-1/6	1/24	
	11/24	-1/2	-1/4	1/2	-1/6	
	11/24	1/2	-1/4	-1/2	1/4	
	1/24	1/6	1/4	1/6	-1/6	
	0	0	0	0	1/24	
6	1/120	-1/24	1/12	-1/12	1/24	-1/120
	13/60	-5/12	1/6	1/6	-1/6	1/24
	11/20	0	-1/2	0	1/4	-1/12
	13/60	1/24	1/6	-1/6	-1/6	1/12
	1/120	1/24	1/12	1/12	1/24	-1/24
	0	0	0	0	0	1/120

B-Spline Piecewise Polynomials

```

PPBS 1 2000 WRITE(6,12010) LFLAG
PPBS 2 FORMAT( ' ? CONVERT ERROR # '15)
PPBS 3 STOP
PPBS 4 END
PPBS 5 C***CONVRT***
PPBS 6 C Computes piecewise polynomial from basis coefficients
PPBS 7 C
PPBS 8 C CALL CONVKT( K, LT,T, NY,LA,A,NA, LCK,LCM,XP,DPP, NXP, LFLAG )
PPBS 9 C
PPBS 10 Reference: J.W. Lewis, M.H. Schultz, S.C. Eisenstat
PPBS 11 "FITS: A Subroutine Package for Spline Regression",
PPBS 12 research report #29, Department of Computer Science,
PPBS 13 Yale University
PPBS 14 C
PPBS 15 C 1) Unless otherwise indicated, variable types are
PPBS 16 C IMPLICIT REAL(A-H, O-Z)
PPBS 17 C IMPLICIT INTEGER(I-N)
PPBS 18 C 2) Value variables (V) pass a value to the subroutine and must
PPBS 19 C have been set by the user or by a previously called subroutine.
PPBS 20 C Result variables (R) return results to the calling subprogram.
PPBS 21 C
PPBS 22 C V K - Degree +1 of spline (less than 10)
PPBS 23 C V T(LT) - B-spline knots
PPBS 24 C V A(LA,NY) - B-spline basis coefficients
PPBS 25 C V NA - Number of B-spline coefficients
PPBS 26 C K XP(LCM+1) - Knots
PPBS 27 C R DPP(LCK,LCM,NY) - Piecewise polynomial stored as the K derivatives
PPBS 28 C of the spline at each knot XP(I) for I=1,NXP
PPBS 29 C R NXP - Number of intervals+1
PPBS 30 C R IFLAG - Status
PPBS 31 C 0 - OK
PPBS 32 C 1 - K GT KMAX (KMAX = 10)
PPBS 33 C 2 - K GT LCK
PPBS 34 C 3 - NXP GT LCM
PPBS 35 C
PPBS 36 C SUBROUTINE CONVKT( K, LT,T, NY,LA,A,NA, LCK,LCM,XP,DPP, NXP,
PPBS 37 C LFLAG )
PPBS 38 C DIMENSION T(LT), A(LA,NY), XP(LCM), DPP(LCK,LCM,NY)
PPBS 39 C The dimensions of the following arrays determine KMAX
PPBS 40 C DATA KMAX/10/
PPBS 41 C KNL = K-1
PPBS 42 C IFLAG = 0
PPBS 43 C IF ( K.GT.LCK ) GO TO 1001
PPBS 44 C IF ( K.GT.KMAX ) GO TO 1002
PPBS 45 C DO 400 JNY=1,NY
PPBS 46 C JR = 0
PPBS 47 C JTO = 0
PPBS 48 C NXP = 0
PPBS 49 C DO 300 JT=K,NA
PPBS 50 C IF ( T(JT).GE.T(JT+1) ) GO TO 300
PPBS 51 C NXP = NXP+1
PPBS 52 C IF ( NXP.GT.LCM ) GO TO 1003
PPBS 53 C XP(NXP) = T(JT)
PPBS 54 C
PPBS 55 C Difference basis coefficients
PPBS 56 C JD = JT-JTO

```

```

PPBS 1 *****
PPBS 2 * Program PPBAS
PPBS 3 * B-spline Piecewise Polynomials
PPBS 4 * 29 Aug 1974
PPBS 5 *****
PPBS 6 EXTERNALS: None
PPBS 7 C
PPBS 8 Reference: S.C. Eisenstat, J.W. Lewis, M.H. Schultz,
PPBS 9 "A Real-Time Algorithm for Least Squares Splines and
PPBS 10 Its Application in Computer Aided Geometric Design",
PPBS 11 research report #29, Department of Computer Science,
PPBS 12 Yale University
PPBS 13 C
PPBS 14 PPBAS is a program to compute the piecewise polynomial
PPBS 15 representations for the B-spline basis over a uniform mesh.
PPBS 16 C
PPBS 17 C Do not change KMAX without changing array dimensions
PPBS 18 C DATA KMAX/10/
PPBS 19 C DATA LI/20/, NY/1/, LA/10/, LCK/10/, LCM/1/
PPBS 20 C DIMENSION A(10), T(20), PP(10), XP(2)
PPBS 21 C
PPBS 22 C Set knots
PPBS 23 C KK = 2*KMAX
PPBS 24 C DO 10 I=1,KK
PPBS 25 C T(I) = FLOMT(I)
PPBS 26 C
PPBS 27 C Set basis coefficients
PPBS 28 C DO 20 I=1,KMAX
PPBS 29 C A(I) = 0.
PPBS 30 C
PPBS 31 C Compute and Print piecewise polynomial
PPBS 32 C
PPBS 33 10010 WRITE(6,10010) KMAX, (I, I=1,KMAX)
PPBS 34 C FORMAT( ' Piecewise Polynomials for B-splines to Order ',
PPBS 35 C 1 I3/' Order
PPBS 36 C Coefficients'/(7X,6I11))
PPBS 37 C Loop through K
PPBS 38 C DO 130 K=1,KMAX
PPBS 39 C WRITE(6,11010) K
PPBS 40 C FORMAT(15)
PPBS 41 C Loop through non-vanishing basis functions
PPBS 42 C DO 120 L=1,K
PPBS 43 C A(L) = 1.
PPBS 44 C CALL CONVKT( K, LT,T, NY,LA,A,K, LCK,LCM,XP,PP, NXP,
PPBS 45 C LFLAG )
PPBS 46 C
PPBS 47 C Convert derivatives at knots to polynomial coefficients
PPBS 48 C IF ( LFLAG.NE.0 ) GO TO 200
PPBS 49 C F = 1.
PPBS 50 C DO 110 I=1,K
PPBS 51 C PP(I) = PP(I)*F
PPBS 52 C F = F/FLOAT(I)
PPBS 53 C
PPBS 54 C Print results
PPBS 55 C WRITE(6,11020) (PP(J), J=1,K)
PPBS 56 C A(L) = 0.
PPBS 57 C 120 CONTINUE
PPBS 58 C 130 CONTINUE
PPBS 59 C 11020 FORMAT(10X,6F11.7)
PPBS 60 C STOP
PPBS 61 C Error in CONVKT

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DO 20 J=L,JD
JP = JR+1
IF (JP.GT.K) JP = 1
JTOJ = JTO + J
AT(L,JP) = A(JTOJ,JNY)
LU = K-JD+J-1
IF (LU.LE.0) GO TO 20
DO 10 LL=1,LU
  KLL = K-LL
  AT(LL+1,JP) = ( AT(LL,JP) - AT(LL,JR) )
  / ( T(JTOJ+KLL) - T(JTOJ) ) * KLL
1
20
30
C Find derivatives at knot
L = K
VB(1) = 1.
DO 220 LK=1,K
LKW1 = LK-1
IF ( LKW1.LE.0 ) GO TO 200
DP(LKW1) = T(JT+LKW1) - T(JT)
DM(LKW1) = T(JT) - T(JT-LKW1+1)
VMP = 0.
DO 110 I=1,LKW1
  VM = VB(I)/(DP(I) + DM(LK-I))
  VBP = VM * DM(LK-I)
  VB(I) = VBP + VMP
  VB(LK) = VMP
110
120
200
C Compute (L-1) derivative at XP(JT) from AT(L,*) and VB(*)
V = 0.
JP = JR
DO 210 J=L,K
  V = V + AT(L,JP)*VB(LK-J+L)
  JP = JP-1
IF (JP.LE.0) JP = K
CONTINUE
EPE(L,NXP,JNY) = V
L = L-1
JIO = JT
300
400
CONTINUE
RETURN
END

C Error returns
1003 LFLAG = LFLAG+1
1002 LFLAG = LFLAG+1
1001 LFLAG = LFLAG+1
RETURN
END

```